

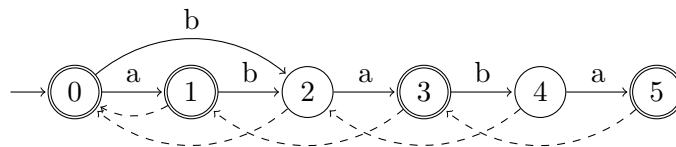
## Tutorial Week 9

**Definition 1.** The suffix automaton corresponding to some text  $y$ , can be regarded as a minimal deterministic final state machine which is a compressed version of a suffix trie. In it, the states are associated to classes representing factors of the string, and their number is upper bounded by  $2n - 1$ , for a text of length  $n$ . The number of arcs, in their turn, are at most  $3n - 4$ . One obtains a suffix automaton with the maximum number of states for a string of the form  $ab^k$ , while for  $ab^k c$  we get an automaton with the maximum number of arcs. The suffix links for a suffix automaton, must not be confused with those of a suffix tree (although they bare the same name). For a state  $p$ , its suffix link points to a different state associated to its longest suffix occurring in a different right context. As for the suffix trees, in the intend to save on space usage, we can merge together all “non-forks”, to obtain a compact suffix automaton.

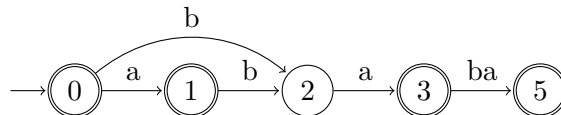
**Exercise 1.** Consider the following list of sequences:  $ababa$ ,  $abcacabb$ ,  $abcacababc$ , and  $abacabacab$ . For each of them, construct the corresponding suffix automaton, including the corresponding suffix links. Finally, construct the associated compact suffix automaton.

What is the complexity of the algorithm? Think about the space usage in each of the cases.

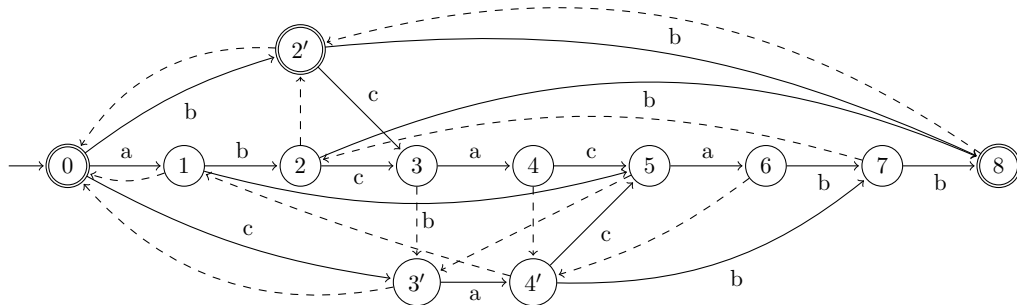
*Solution:* The suffix automaton associated to  $ababa$  is



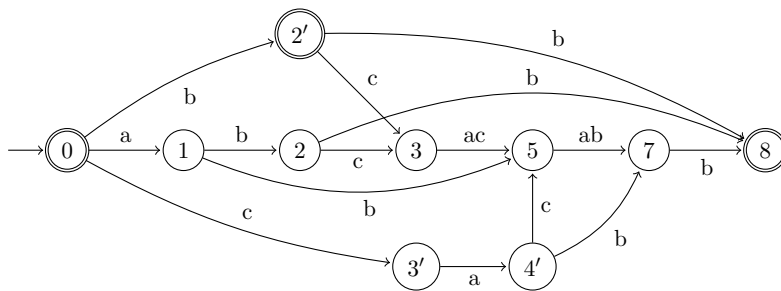
The CSA corresponding to this automaton is



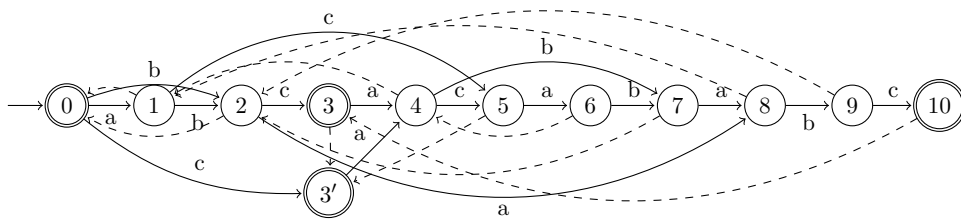
The suffix automaton associated to  $abcacabb$  is



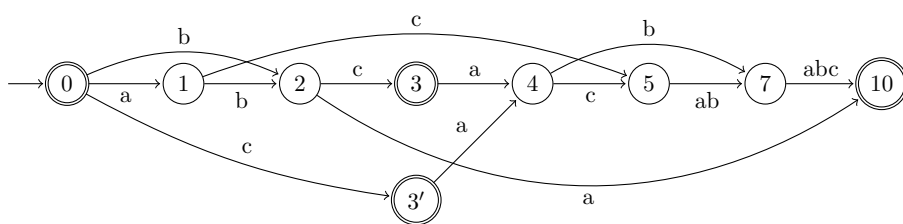
The *CSA* corresponding to this automaton is



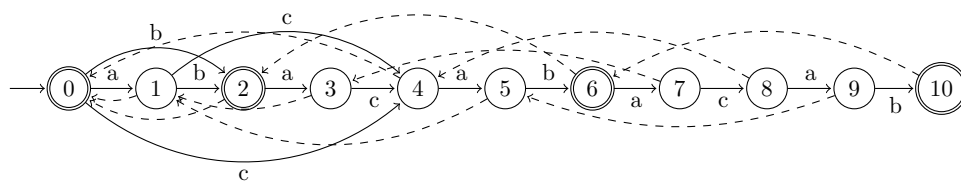
The suffix automaton associated to *abcacababc* is



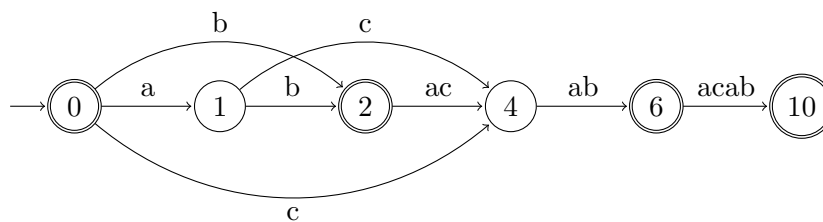
The *CSA* corresponding to this automaton is



The suffix automaton associated to *abacabacab* is



The *CSA* corresponding to this automaton is



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