

Tutorial Week 7

Definition 1. The suffix table SUF of a string y of length n gives us the order of the suffixes of y , such that $\text{SUF}[i] < \text{SUF}[i + 1]$. The LCP array corresponding to y gives us the lengths of the longest common prefixes of consecutive elements in the suffix table:

$$\text{LCP}[i] = |\text{lcp}(y[\text{SUF}[i - 1]..n - 1], y[\text{SUF}[i]..n - 1])|.$$

Together the two form the suffix array. If by $\text{First}_k(u)$ we denote the prefix of length k of u , when $|u| \geq k$, and u , otherwise, then by $R_k[i]$ we denote the rank of $\text{First}_k(y[i..n - 1])$ inside the ordered list of all $\text{First}_k(u)$, where u is a non-empty suffix of y .

Exercise 1. Consider the following list of sequences: *ababa*, *abcacabb*, *abcacababc*, and *abacabacab*. For each of them, using the prefix doubling technique find out the corresponding ranks. Next, find out their corresponding suffix array consisting of the suffix table SUF and the corresponding LCP table, respectively.

What is the complexity of the algorithm?

Solution:

i/r	0	1	2	3	4
$y[i]$	<i>a</i>	<i>b</i>	<i>a</i>	<i>b</i>	<i>a</i>
$R_1[i]$	0	1	0	1	0
$R_2[i]$	1	2	1	2	0
$R_4[i]$	2	4	1	3	0
$\text{SUF}[r]$	4	2	0	3	1
$\text{LCP}[r]$	0	1	3	0	2

i/r	0	1	2	3	4	5	6	7
$y[i]$	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>b</i>
$R_1[i]$	0	1	2	0	2	0	1	1
$R_2[i]$	0	4	5	1	5	0	3	2
$R_4[i]$	1	5	7	2	6	0	4	3
$\text{SUF}[r]$	5	0	3	7	6	1	4	2
$\text{LCP}[r]$	0	2	1	0	1	1	0	2

i/r	0	1	2	3	4	5	6	7	8	9
$y[i]$	a	b	c	a	c	a	b	a	b	c
$R_1[i]$	0	1	2	0	2	0	1	0	1	2
$R_2[i]$	0	3	5	1	5	0	2	0	3	4
$R_4[i]$	2	6	9	3	8	0	4	1	5	7
$\text{SUF}[r]$	5	7	0	3	6	8	1	9	4	2
$\text{LCP}[r]$	0	2	3	1	0	1	2	0	1	2

i/r	0	1	2	3	4	5	6	7	8	9
$y[i]$	a	b	a	c	a	b	a	c	a	b
$R_1[i]$	0	1	0	2	0	1	0	2	0	1
$R_2[i]$	0	3	1	4	0	3	1	4	0	2
$R_4[i]$	1	4	2	6	1	4	2	5	0	3
$R_8[i]$	2	7	4	9	1	6	3	8	0	5
$\text{SUF}[r]$	8	4	0	6	2	9	5	1	7	3
$\text{LCP}[r]$	0	2	6	1	4	0	1	5	0	3

The algorithm takes linear time to find the array SUF using the skew algorithm and linear time to compute the array LCP . Please note that when computing the $\text{SUF}[r]$ we look at the positions of the ranks in the word! Furthermore, when computing the $\text{LCP}[r]$ we look at the $r - 1$ and r positions in the array SUF and not in the word! ■

Exercise 2. Consider now the same list of sequences: *ababa*, *abcacabb*, *abcacababc*, and *abacabacab*. For each of them, using the skew algorithm find out the corresponding ranks, considering the t and s sequences. Next, find out their corresponding suffix array consisting of the suffix table SUF and the corresponding LCP table (obviously, these are the same as for the previous exercise).

What is the complexity of the algorithm?

Solution: First we look at the string *ababa*.

i	0	1	2	3	4
$y[i]$	a	b	a	b	a

Rank t		Rank s	i	Suff(1230)	Rank	j	$(y[j], s[j+1])$
0	a	0	4	0	0	2	$(a, 2)$
1	aba	1	0	1230			
2	ba	2	3	230			
3	bab	3	1	30			

i	0	1	2	3	4
$y[i]$	a	b	a	b	a
$r[i]$	2	4	1	3	0
SUF[r]	4	2	0	3	1
LCP[r]	0	1	3	0	2

Next we look at the string $abcacabb$.

i	0	1	2	3	4	5	6	7
$y[i]$	a	b	c	a	c	a	b	b

Rank t		Rank s	i	Suff(013452)	Rank	j	$(y[j], s[j+1])$
0	abc	0	0	013452	0	2	$(a, 3)$
1	aca	1	3	13452	1	5	$(c, 1)$
2	b	2	7	2			
3	bb	3	6	3452			
4	bca	4	1	452			
5	cab	5	4	52			

i	0	1	2	3	4	5	6	7
$y[i]$	a	b	c	a	c	a	b	b
$r[i]$	1	5	7	2	6	0	4	3
SUF[r]	5	0	3	7	6	1	4	2
LCP[r]	0	2	1	0	1	1	0	2

Next we look at the string $abcacababc$.

i	0	1	2	3	4	5	6	7	8	9
$y[i]$	a	b	c	a	c	a	b	a	b	c

Rank t		Rank s	i	Suff(0124350)	Rank	j	$(y[j], s[j+1])$
0	<i>abc</i>	0	7	0	0	5	$(a, 2)$
1	<i>aca</i>	1	0	0124350	1	8	$(b, 4)$
2	<i>bab</i>	2	3	124350	2	2	$(c, 1)$
3	<i>bca</i>	3	6	24350			
4	<i>c</i>	4	1	350			
5	<i>cab</i>	5	9	4350			
		6	4	50			

i	0	1	2	3	4	5	6	7	8	9
$y[i]$	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>a</i>	<i>b</i>	<i>c</i>
$r[i]$	2	6	9	3	8	0	4	1	5	7
SUF[r]	5	7	0	3	6	8	1	9	4	2
LCP[r]	0	2	3	1	0	1	2	0	1	2

Finally, we look at the string *abacabacab*.

i	0	1	2	3	4	5	6	7	8	9
$y[i]$	<i>a</i>	<i>b</i>	<i>a</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>a</i>	<i>c</i>	<i>a</i>	<i>b</i>

Rank t		Rank s	i	Suff(0412304)	Rank	j	$(y[j], s[j+1])$
0	<i>aba</i>	0	4	04	0	8	$(a, 2)$
1	<i>aca</i>	1	0	0412304	1	2	$(a, 4)$
2	<i>b</i>	2	6	12304	2	5	$(b, 1)$
3	<i>bac</i>	3	9	2304			
4	<i>cab</i>	4	1	304			
		5	7	4			
		5	3	412304			

i	0	1	2	3	4	5	6	7	8	9
$y[i]$	<i>a</i>	<i>b</i>	<i>a</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>a</i>	<i>c</i>	<i>a</i>	<i>b</i>
$r[i]$	2	7	4	9	1	6	3	8	0	5
SUF[r]	8	4	0	6	2	9	5	1	7	3
LCP[r]	0	2	6	1	4	0	1	5	0	3

The algorithm takes linear time to find the array SUF using the skew algorithm and linear time to compute the array LCP. Please note that when computing the SUF[r] we look at the positions of the ranks in the word! Furthermore, when computing the LCP[r] we look at the $r - 1$ and r positions in the array SUF and not in the word! ■