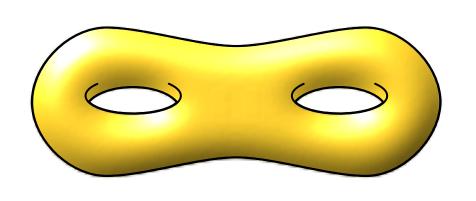
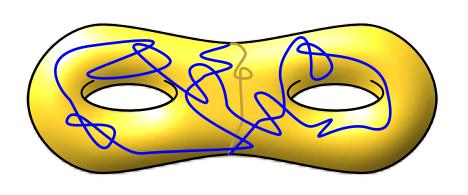
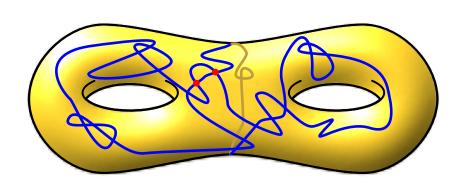
Autour du nombre géométrique d'intersection

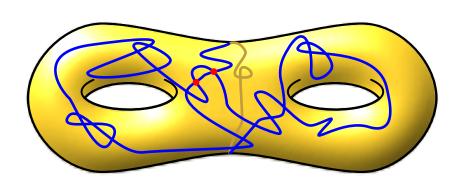
Francis Lazarus
GIPSA-Lab, Grenoble, CNRS

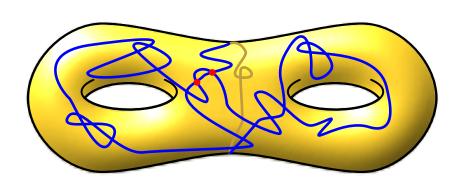


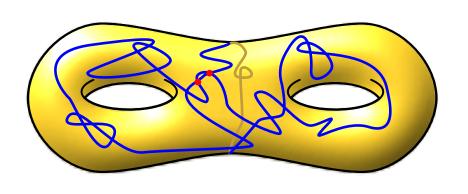


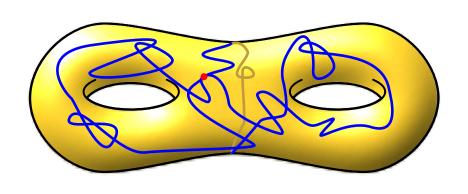


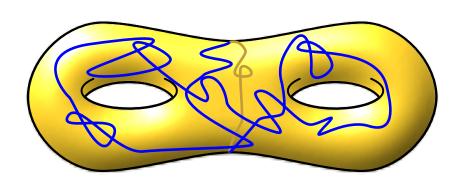


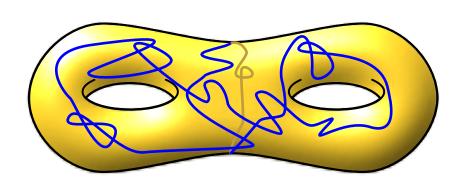


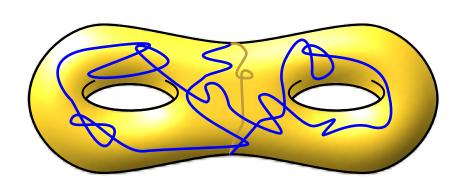


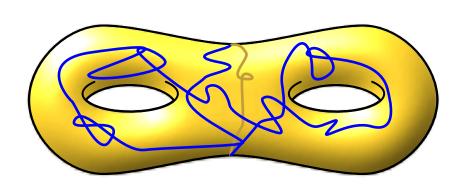


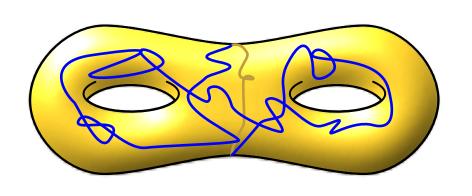


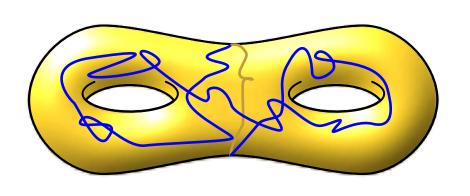


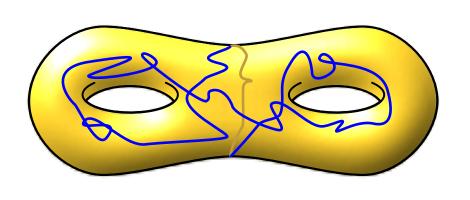


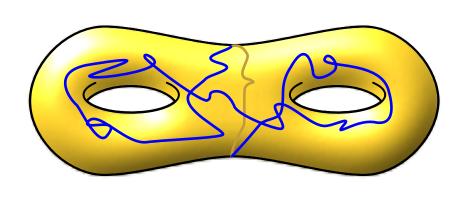


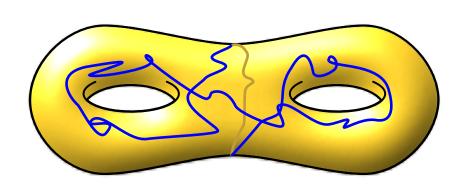


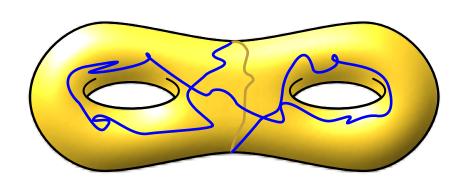


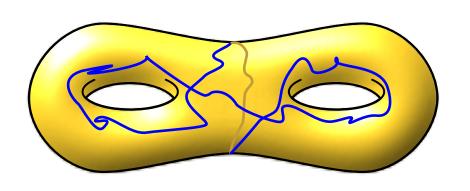


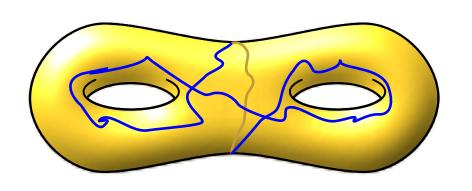


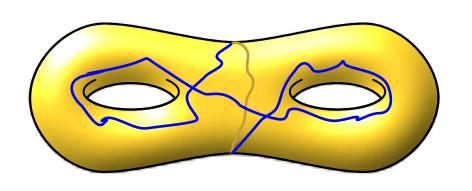


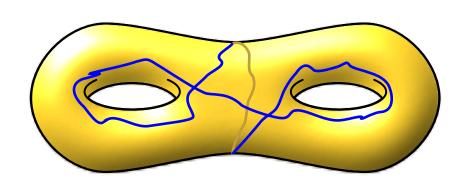


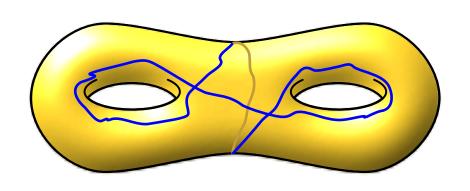


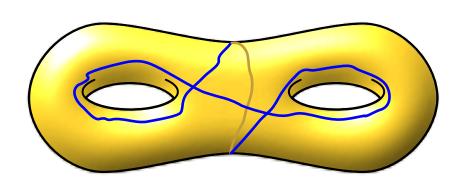


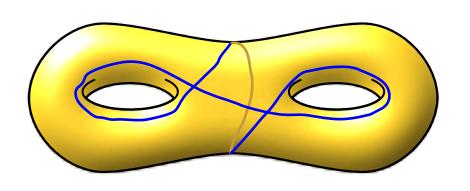


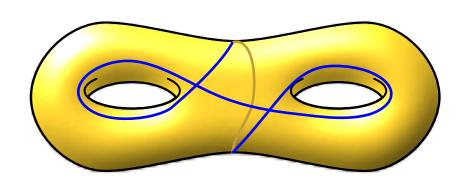


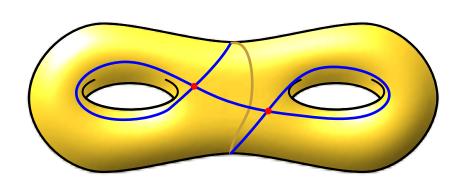


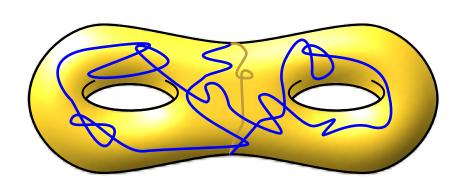


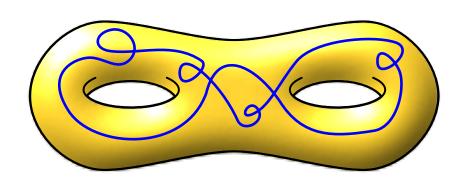


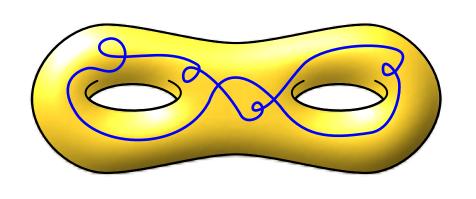


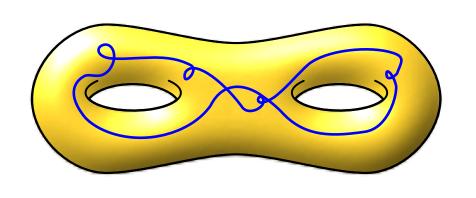


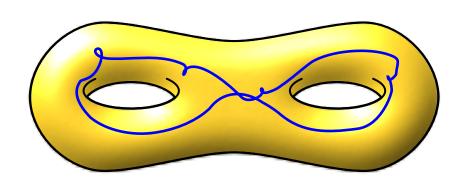


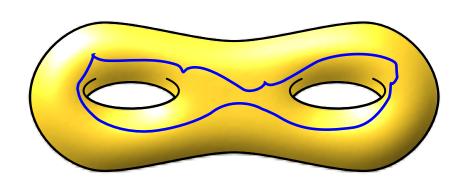


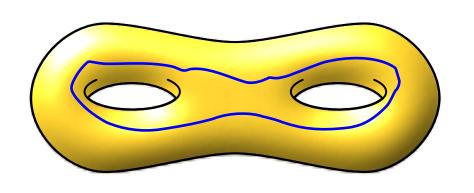


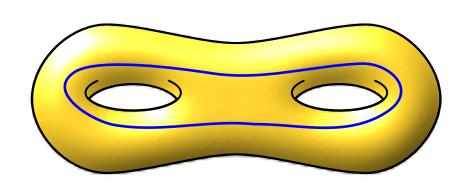


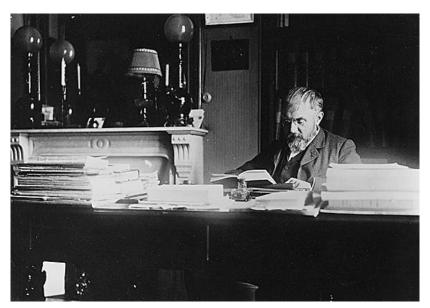












1854 - 1912

ANALYSIS SITUS

Journal de P.Émile Polytechnique, 5-5, p. 1-121 (1965)

Introduction.

La Géométrie à a dimensiona a un objet réel; personne n'en deute najeccd'uni. Les étens de l'Appeneupes must unespillules de délicitions précluse cousses ceux de l'espece coétaines, et si nous se provente cous les représentates; nous pouvent les causevrier et les duilles. Si danc, par emangla, la Méxanigne a plas de testis dimensions dels mos condamante cousse dépouveue de son abjet. Il viva en cas de traits de l'Hémonétories.

La Géométrie, en effic, n'a pas pour unique raison d'être la description immédiate des corps qui toublest suu nou seus : élle est sumit test l'étude saulptique d'un groupe; riun n'empéche, par consèquent, d'aborder d'entres gourges saulegres et plus georieuxe.

Mais passquai, dista-bon, no pas construe la langua analytique et le resplace par sa lacqua geneticipes, qui perf con sea sunauger dat que la sesta no persona plus interprete. Cest que ce langua excevar set plus concisi a éven consiste que l'analogie sere la Géométric celinaire pest ceter des succisions d'Alga Respués y sungitere des plustralisations sollar.

Pent-tier can minum un secot-alles pas seffimates? Ce n'ost pas naues, un effet, qu'une solucce nois législane : il faut que l'utilisé un pointe en être contretée. Tent d'objets divers sollicitant notre attention, que los plus importants ent souls d'evit de l'obtenir.

SUR

CERTAINES SURFACES ALGÉBRIQUES; TROISIÈME COMPLÉMENT

A L'ANALYSIS SITUS

mateix de la factiof trathématique de France, t. Ill. n. (a-re-) (ann).

Proposiza-aque d'étadier, su point de vas de l'Analyste attur, la eurhee

(i) $x = \sqrt{F(x,y)},$ of F(x,y) as in polynome. Nous suppose one que la courbe

F(s,y)=s so présent que des poists ordinaires et $\frac{d}{ds}$ n'est pas uni, so bien et $\frac{d\theta}{ds}=s$, mais sem que $\frac{d^2p}{ds} + \frac{d\theta}{ds} \frac{ds}{ds} + s$, cade poists dyables ordinaires et $\frac{d\theta}{ds} = s$.

$$\frac{dT}{dz} = \frac{dT}{dy} = 0,$$

nais saux que $\frac{\partial Z}{\partial z}$ y $\frac{\partial Z}{\partial z}$ $\frac{\partial Z}{\partial z}$ $\frac{\partial Z}{\partial z}$ ($\frac{\partial Z}{\partial z}$) virtualent. Considèrant d'Aberl y comme sau constatue. Aborr l'équation (y) représeurs une souvée aplitéque et la confunction z et z 'ha point de cette equite-permut l'exprisse, comme en le suit, comme de foncțion fichielente d'em anteu viriale studiire z. Confidente le group franție, corresponddențe le ydyppe dubites qu'il regardire. En getardi, le pilypone fechies qu'i correspond to tu a conché gir game z en troplytone fechies.

COMPLEMENT A L'ANALYSIS SITUS

Anadirani del Circlo Matematico di Pelaresa, s. 15, p. 181-16 (1898).

(i. - Introduction.

Dana le Journal de l'Écule Polytechnique (volume du centennie de la fondetion de l'Écule, 1864). Pai publis en Manules instablé Anadysi nius, où Fristiè les avaitées de l'espoce à plus de trois disconsissa et les propriétés des nombres de Betti. C'est à ou Manules que supporterent les remois que le serie insend à faire fréquentation dans la soite, en mémbronaux avalences le lettre desdréss district desdréss district desdréss district.

Daza ce Minicire su trouve étocock le théorèum utivant : Pour équie variété fermée, les mondres de Batti également distants des aztrémes nont égaux. Le mine théorèum a été étancé par M. Picarl dem su Théorie des fonctions alrebricant de deux variables.

M. Beegard Visit de revenir aur en nebus problème dans un travail tels emanquelle, public en langue dansies, son la titra «Foretadier eil en tapologist keori fer de algebranks Sommenhaps (Coppshages, det Neue Forleg Brust Bejenen, 1898). D'apols lei la théorème en question est inemet el les démancacions son aven valent.

Avent d'exeminer les objections de M. Heeganod, il convient de feire une distinction. Il y a desa manières de définir les nombres de Betti.

Canadánna nos unitát V que je seponenie, par example, formás judint $v_0, v_1, v_1, \dots, v_n$ a unitát à p dimensions, fainat partie de V. Le suppose q d'un no pásico par tourant de unitát à p+1 dimensions faiquat partie de V et fonto, v_1, \dots, v_n canadóns la frendêne complète; mais que, si en loca diplotta un $(a+1)^{m_1}$ unitát de p dissentions que l'populari v_{m_1} au qui fon

SUB.

LES CYCLES DES SURFACES ALGÉBRIQUES:

QUATRIÈME COMPLÉMENT A L'ANALYSIS SITUS

s 1. Introduction

Les benns tresques de M. Friend our les Touytons elgérichypeur est in pagin laspingues que deletare Timpertane de trais des explete à une, desse as syiri dimensions. Dit poste qu'en prometa popilique « entre questies a les présiges que l'il explant desse Les Antonio de une des deux personnes des traises de l'Endoir Polymentogres, trais de constantir; Rendiment des Cerus Mansacand de Palerran, XIII Proceedinges qu'en la containe relation de Cerus Mansacand Solory, « de XXXIII) es l'il viteurs dans cerus restaurs Mathematica Solory, « de XXXIII) es l'il viteurs dans cerus restaurs de l'autre de containe principation de M. Flanck.

Le rappelle qu'etrat denné une varieté V formée à p dimensione, je trace sur ceme varieté d'autou varietés, ferrates en non, d'un moins grand nombre de dimensions; je désigne par W_q une vérieté à q dimensions tracée de la serie sur V.

Si ΣW_p set un ensemble de varients à g dimensions et ΣW_{p-1} un ensemble de varients à g-1 dimensions, la congruence ΣW_{p-1} ΣW_{p-1}

signific (per définition) que ZW_{q-4} foreze la frontière complète de l'ensemble

SECOND COMPLÉMENT

L'ANALYSIS SITUS

Proceedings of the London Mathematical Scotter, 1, 22, p. 111-248 (15 juin 1911).

Televitoritor

Pai public dans le Journal de l'Étade Polymotholyan un travail intitule de destipais rittes y ja et a sols accorpt une notacide fait de même problème dans un Mentrius postus pour tiere e Compférence à l'Andreya ritters, et qui et été imprise dans les Rendiquent del Clerche Metresaciée de Publicars. Organdant le question en les mêtres équains, es je sons une doute farre le compressa de publica en les mêtres équains, es je sons une doute faire de présent le publicar requises. Pour como dois je une homanni le cercitaine canadidavision en si en el acustre à timilitée à déclaire et a consoluir les canadidacitions en les estaturs à timilitée à déclaire et a consoluir les canadis-

Les reuvries perteux simplement une indication de presgraphe ou de page en supporterent en poemier Mémoire, coloi da Journal de l'École Polytemplyer; les reuvris où cre indications servou précôdice de la louve o es supporterent en Minaire des Rendijonni.

poécédemment sequis.

Quantung recoveis aux paragraphes du présent Métacies, je les ferei précéder

CINQUIÈME COMPLÉMENT

L'ANALYSIS SITTIS

Anadiensi del Cirette matematico di Pedermo, I. (6, p. di re (1996).

1. Psi digi co sassemi l'occasion de m'occape d'Analysis situs; j'si d'aberd politic on Manciero sur es signi dons le tuns da Centenatire da Jurmal de l'Elbert Polyprochago; co Mansiero si dun de quetre compléannes qui sus pere jaus le voue XIII des Reculionats del Circilo Matematico di Polermo, dans le voue XXXVII des Proceedings et de London Mathematica Society, dans le Dalieloi de la Society Mathematique de Prince en 1991, et efficie.

dans le Aurenol de Livarille en 1901. Je revieus encore sujuard'hui sur esta mense questica, pervandé qu'en n'en pourre vair à bout que par des effects répétés et qu'elle est asses importante mare les anchies.

Cette fois je me suis horné à l'étude de ceruines variétés à trois dimensions, mais les catcheles que f'ui exployées pourrant être une doute d'un suge plus général. En passant je me suis étende aussi longuement eur certaines posjréétés des courbes forméte une l'on nout invers ser les surfaces ferméte de l'estende des courbes forméte une l'on nout invers ser les surfaces ferméte de l'estende

continuer. Le résolut fisal que j'erais en von est le selvant. Dans le second complément j'ai moutré que pour cesusteiner une varieté il ne suffit pas de commêtre ses namères de Butil, mais que certains coefficients que j'ai appelés coefficients de tonign (accord complément, 8.5, p. 565) issuéent ou relé injussement.

On pourreit alors se demagder si la considération de ces coefficients suffit;

Nous sommes donc conduit à une contradiction, ce qui veut dire que l'hypothèse faite au début était absurde et que le cycle K doit être bouclé.

C. Q. F. D.

4. Nous avons vu au paragraphe précédent qu'il est relativement aisé de reconnaître si un cycle donné est homologue à un cycle non bouclé, ou si deux cycles donnés sont respectivement homologues à deux cycles qui ne se coupent pas. Nous allons dans le présent paragraphe examiner une question analogue :

Comment reconnaître si un cycle donné est équivalent à un cycle non bouclé, ou si deux cycles donnés sont équivalents à deux cycles qui ne se coupent pas?

Mais avant d'aborder cette question, revenons sur la définition de l'équivalence.

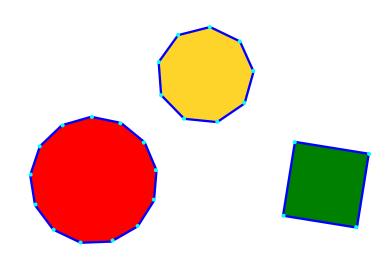
Jusqu'ici nous avons toujours entendu cette équivalence de la façon suivante :

Quand nous écrivons

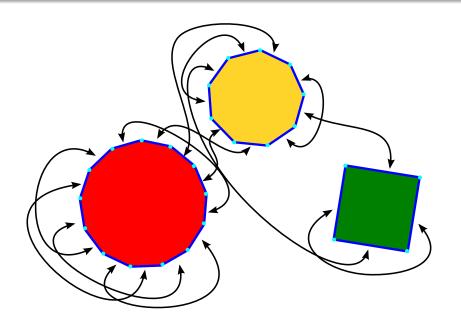
 $C \Longrightarrow C'$

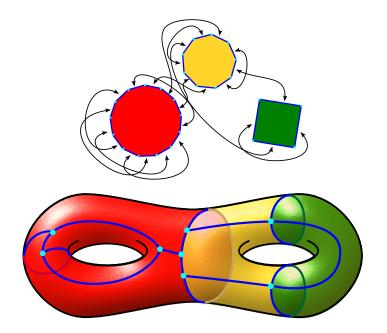
H. P. - VI.

What is a surface?



What is a surface?





Über den Begriff der Riemannschen Fläche.

Von Tibor Radó in Szeged,

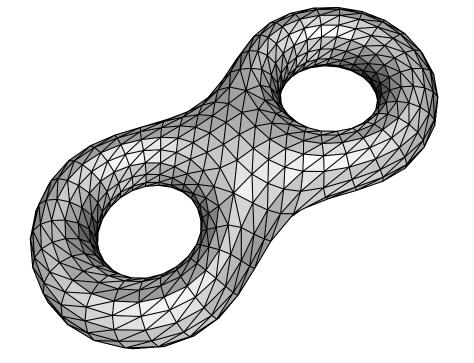
Einleitung.

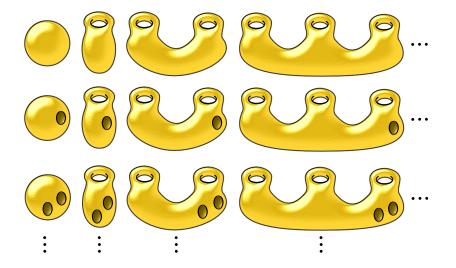
Die vorliegende Arbeit enthält eine Untersuchung, zu welcher ich beim Studium des grundlegenden Werkes des Herrn Wexu. Diber Die Idee der Riemannschen Fläche geführt wurde. Bekanntlich wird in diesem Buche der Begriff der Riemannschen Fläche zum ersten Male in vollkommen strenger Weise erklärt, und zwar wie folgt. Eine Riemannschen Fläche ist eine zweidimensionale Mannigfaltigkeit, welche trianguliert werden kann, und für welche eine konforme Abbildung im Kleinen mitgegeben ist. Wir werden diese Begriffe eingehend besprechen (§§ 3 und 4), müssen aber gleich an dieser Stelle die Forderung der Triangulierbarkeit doch etwas genauer betrachten, um unser Problem formulieren zu können.

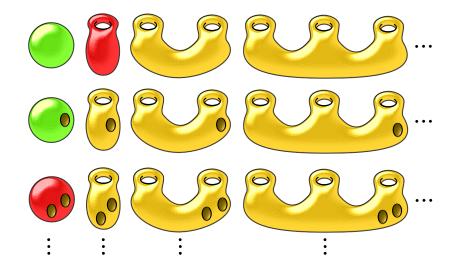
Der Ausdruck zwidimensionale Mannigfaltigkeit wird nicht von allen Autoren in demselben Sinne gebraucht. Sie ist jedenfalls ein zusammenhängender topologischer Raum im HAUSDORFF-schen Sinne, welcher im Kleinen der xy-Ebene hombomorph ist aber es wird manchmal noch die Forderung an sie gestellt, sie soll dem zweiten HAUSDORFF-schen Abzählbarkeitsaxiom genügen. Wollte man den Ausdruck zweidimensionale Mannigfaltigkeit bei der Erklärung der REMANNSChen Fläche in diesem Sinne verstehen, so wäre die Forderung der Triangulierbarkeit überflüssig. Unter Voraussetzung dieses Abzählbarkeitsaxioms bietet nämlich die Triangulierung einer zweidimensionalen Mannigfaltigkeit keine prinzipielle, sondern nur technische Schwierigkeiten, und die ex plicite Forderung der Triangulierbarkeit würde einfach bedeuten dass man mit möglicherweise umständlichen, aber im Grunde ganz einfachen Betrachtungen keine Zeit verlieren will.



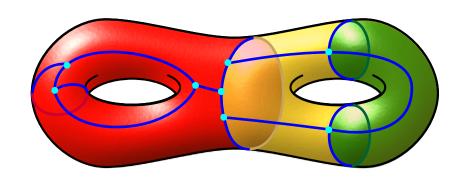
Tibor Radó 1895 – 1965

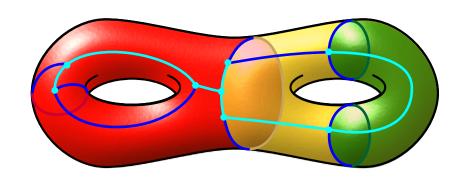


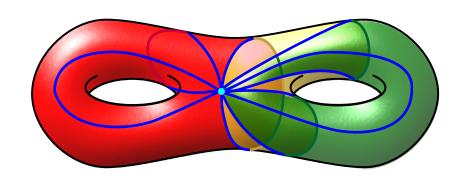


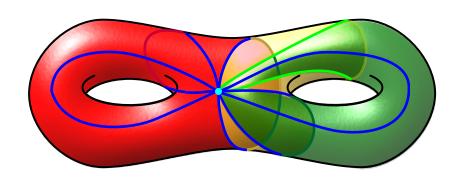


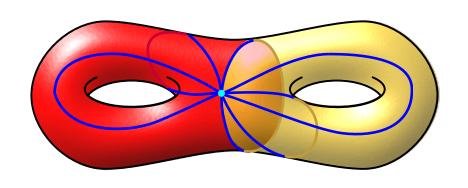
$$S-A+F=\chi=2-2g-b$$

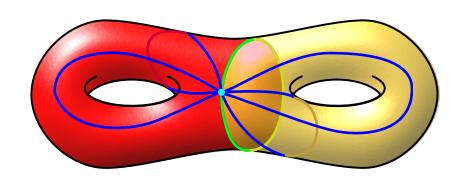


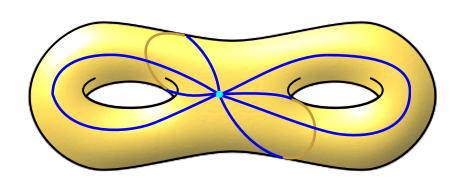


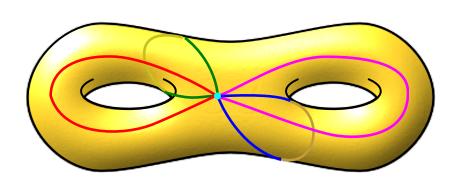


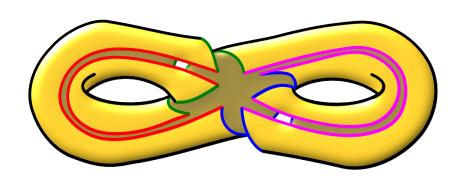


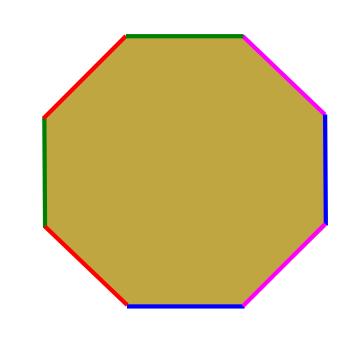




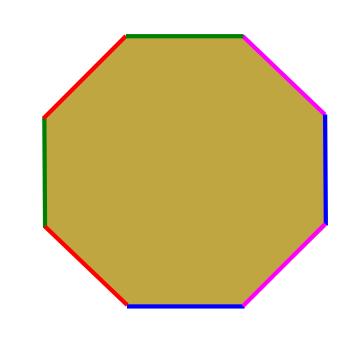


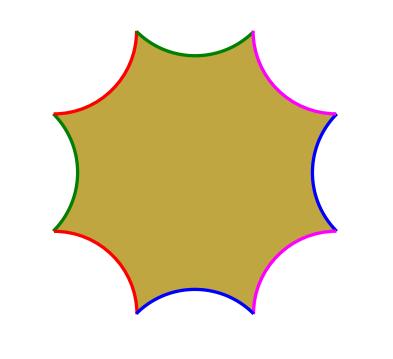






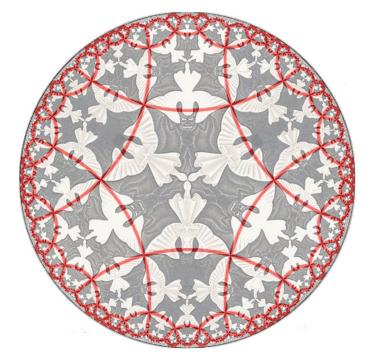
On peut se placer, dans l'étude de la question qui nous occupe, à plusieurs points de vue différents. Représentons d'abord notre surface par un polygone fuchsien R₀ de la première famille, construisons les différents transformés de ce polygone par les transformations du groupe fuchsien correspondant G; ces transformés rempliront le cercle fondamental. Un cycle quelconque C sera alors representé par un arc de courbe MM', allant d'un point M à un de ses transformés M'. Deux cycles proprement équivalents seront représentés par deux arcs de courbe MPM' et MQM' ayant mêmes extrémités et réciproque-

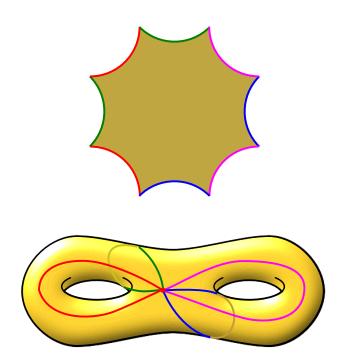


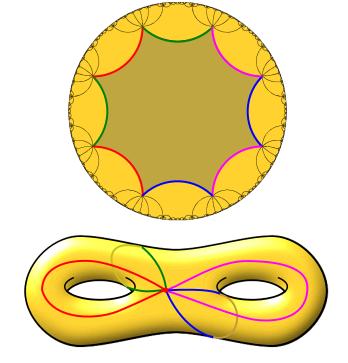


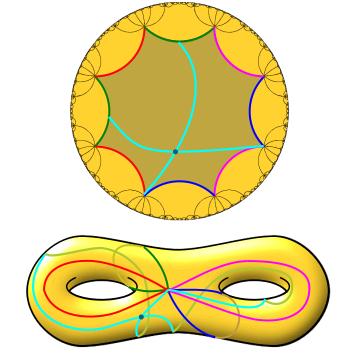


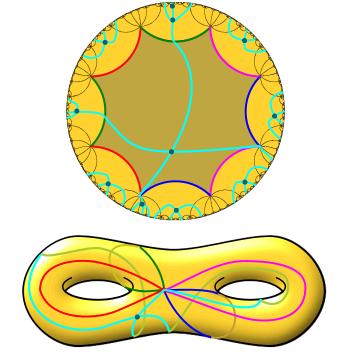


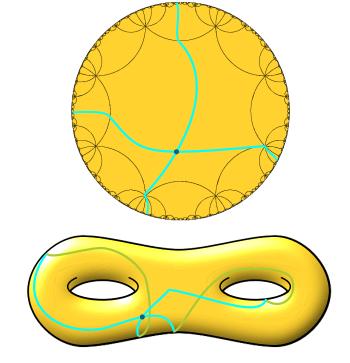


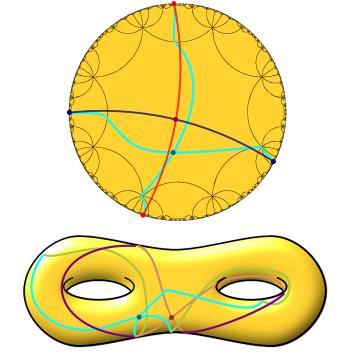


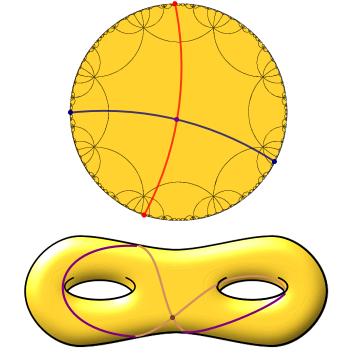












ALGORITHMS FOR JORDAN CURVES ON COMPACT SURFACES

By Bruce L. Reinhart* (Received October 20, 1960)

As is well known, free homotopy classes of mappings of the circle into a space correspond to conjugacy classes in the fundamental group. If the space is two-dimensional, such a free homotopy class may or may not contain a simple closed (that is, Jordan) curve. Our purpose here is to give an algorithm for determining which free homotopy classes admit such a curve in the case of compact surfaces of negative Euler number. Our algorithm is applicable to orientable and non-orientable surfaces, with or without boundary. This problem was first discussed by Dehn [4] and later by Baer [1] and Goeritz [5]. All of these studies are based on cutting the surface into spheres with 3 holes and deriving a (hopefully) canonical form for each free homotopy class by pasting together curves on each sphere with holes. Except for the surface of genus 2 without boundary, no complete answer has been achieved in this way. We shall treat the problem globally by imbedding the fundamental group into the group of motions of the hyperbolic plane, in the spirit of Poincaré [14] and numerous works of Nielsen. As a preliminary, we give in §1 a simple method for assigning to each word in the usual generators of the fundamental group a curve on the surface which has double points only in the neighborhood of the base point, and no other multiple points. We call such a curve an indicating curve for the word. In § 2, we apply the fact that each motion of the hyperbolic plane which arises in our problem leaves fixed a unique geodesic, its axis. The projection of an axis onto the surface is the unique geodesic lying in the free homotopy class containing the motion of which it is axis. For orientable surfaces, a free homotopy class admits a simple closed curve if and only if its geodesic is simple [14, p. 467], while

- Henri Poincaré. V-ème complément à l'analysis situs. Rendiconti del Circolo Matematico di Palermo, 18(1):45 –110, 1904.
- Bruce L. Reinhart. Algorithms for jordan curves on compact surfaces.
 Ann. of Math., p. 209–222, 1962.
- Heiner Zieschang. Algorithmen für einfache kurven auf flächen. Mathematica Scandinavica, 17:17–40, 1965.
- Heiner Zieschang. Algorithmen für einfache kurven auf flächen II.
 Mathematica scandinavica, 25:49–58, 1969.
- David R.J. Chillingworth. Simple closed curves on surfaces. Bull. of London Math. Soc., 1(3):310–314, 1969.
- London Math. Soc., 1(3):310–314, 1969.
 David R.J. Chillingworth. An algorithm for families of disjoint simple closed curves on surfaces. Bull. of London Math. Soc., 3(1):23–26,

1971.

- David R.J. Chillingworth. Winding numbers on surfaces. II.
 Mathematische Annalen, 199(3):131–153, 1972.

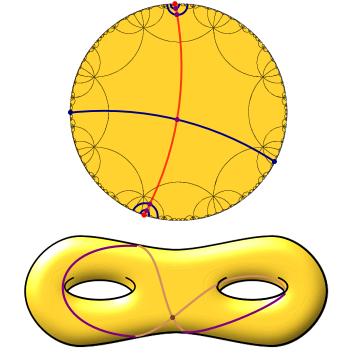
 Joan S. Birman and Caroline Series. An algorithm for simple curves on.
- Joan S. Birman and Caroline Series. An algorithm for simple curves on surfaces. J. London Math. Soc., 29(2):331–342, 1984.

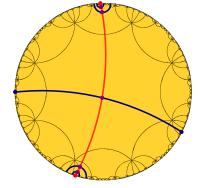
- Marshall Cohen and Martin Lustig. Paths of geodesics and geometric intersection numbers: I. Ann. of Math. Stud. 111:479–500, 1987.
- Martin Lustig. Paths of geodesics and geometric intersection numbers: II. *Ann. of Math. Stud.* 111:501–543, 1987.
- Joel Hass and Peter Scott. Shortening curves on surfaces. *Topology*, 33(1):25–43, 1994.
 Joel Hass and Peter Scott. Configurations of curves and geodesics on
- surfaces. Geometry and Topology Monographs, 2:201–213, 1999.
 Maurits de Graaf and Alexander Schrijver. Making curves minimally crossing by Reidemeister moves. J. Combinatorial Theory, Series B.
- Max Neumann-Coto. A characterization of shortest geodesics on surfaces. Algebraic and Geometric Topology, 1:349–368, 2001.

70(1):134–156, 1997.

- J.M. Paterson. A combinatorial algorithm for immersed loops in surfaces. *Topology and its Applications*, 123(2):205–234, 2002.
- Daciberg L. Gonçalves, Elena Kudryavtseva, and Heiner Zieschang. An algorithm for minimal number of (self-)intersection points of curves on surfaces. *Proc. of the Seminar on Vector and Tensor Analysis*, 26:139–167, 2005.

	closed		free	special
	surface	counting	homotopy	feature
Chillingworth				winding
'69				number
Birman &				retraction
Series '84			√	onto a graph
Cohen &				retraction
Lustig '87		\checkmark	√	onto a graph
				canonical
Lustig '87	\checkmark	$\overline{}$	√	representative
de Graaf &				Reidemeister
Schrijver '97	\checkmark	$\overline{}$	√	moves
				Reidemeister
Paterson '02	\checkmark	$\overline{}$	$\overline{}$	moves
Gonçalves	_	_	_	algebraic
et al. '05	\checkmark		√	approach





- If a curve c is primitive its lifts are uniquely defined by their limit points.
- If au is the hyperbolic translation corresponding to a lift $ilde{c}_0$ of a primitive c then

 $\imath(c) = |\{ ext{set of pairs of limit points crossing } ilde{c}_0 \} / \langle au
angle | \}$

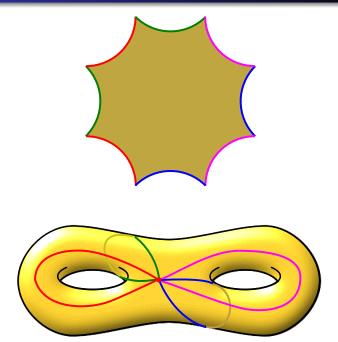
- If a curve c is primitive its lifts are uniquely defined by their limit points.
- If au is the hyperbolic translation corresponding to a lift $ilde{c}_0$ of a primitive c then

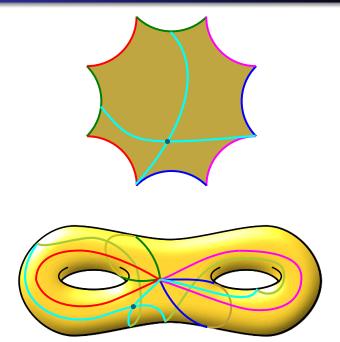
 $\imath(c) = |\{ ext{set of pairs of limit points crossing } \tilde{c}_0 \} / \langle au
angle |$

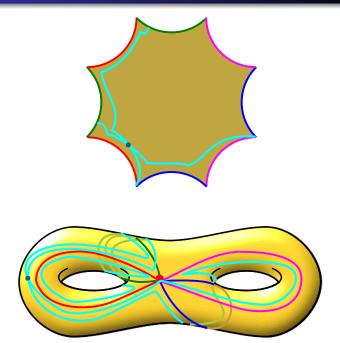
The plan

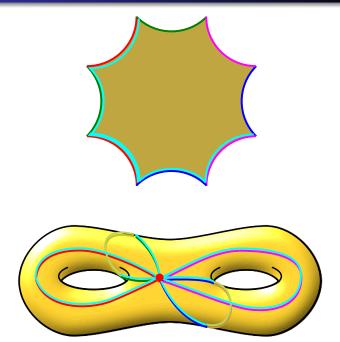
For a given curve c

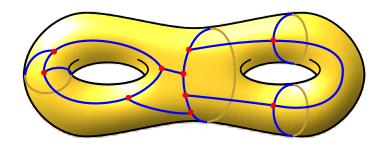
- Determine the primitive root of c.
- 2 Count the number of classes of crossing pairs of limit points (for the root of *c*).
- Use adequate formula if c is not primitive.

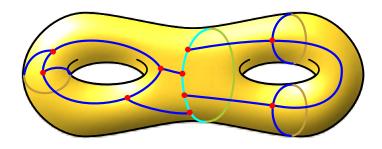


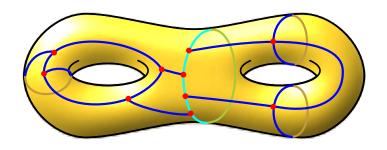


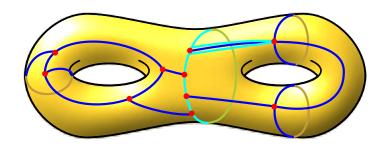






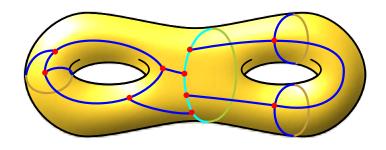






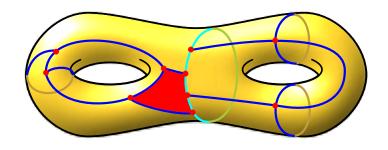
Two rules:

• add/delete a spur along an edge

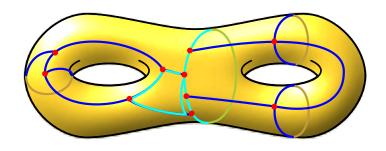


Two rules:

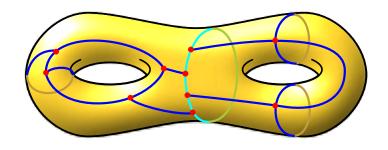
add/delete a spur along an edge



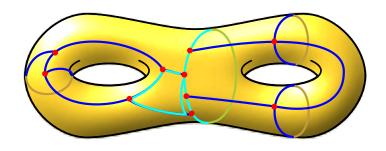
- add/delete a spur along an edge
- add/delete a facial walk



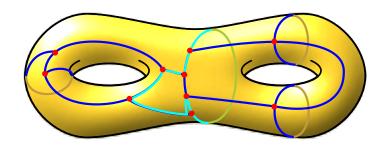
- add/delete a spur along an edge
- add/delete a facial walk



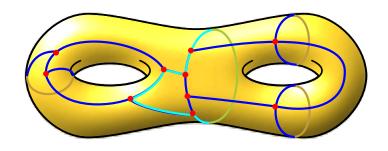
- add/delete a spur along an edge
- add/delete a facial walk



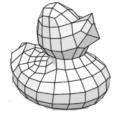
- add/delete a spur along an edge
- add/delete a facial walk

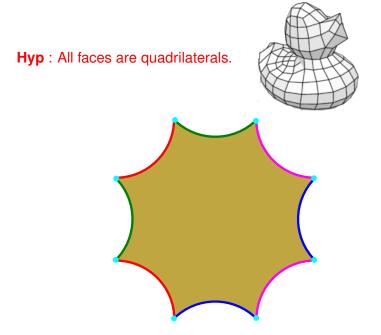


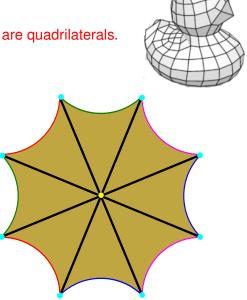
- add/delete a spur along an edge
- add/delete a facial walk

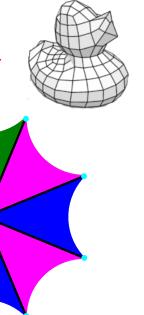


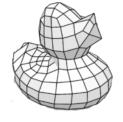
- add/delete a spur along an edge
- add/delete a facial walk







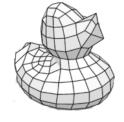




Discrete curvature

$$\kappa_s = 1 - \frac{d_s}{2} + \frac{c_s}{4}$$

 $d_s :=$ degree of s and $c_s :=$ number of incident corners.



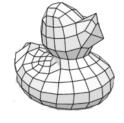
Discrete curvature

$$\kappa_s = 1 - \frac{d_s}{2} + \frac{c_s}{4}$$

 $d_s :=$ degree of s and $c_s :=$ number of incident corners.

$$\sum_{\mathbf{s}\in\mathcal{S}}\kappa_{\mathbf{s}}=\chi$$

PROOF.
$$\sum_{s \in S} \kappa_s =$$



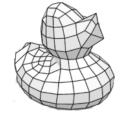
Discrete curvature

$$\kappa_s = 1 - \frac{d_s}{2} + \frac{c_s}{4}$$

 $d_s :=$ degree of s and $c_s :=$ number of incident corners.

$$\sum_{\mathbf{s}\in\mathcal{S}}\kappa_{\mathbf{s}}=\chi$$

PROOF.
$$\sum_{s \in S} \kappa_s = S -$$



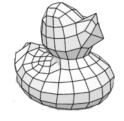
Discrete curvature

$$\kappa_s = 1 - \frac{d_s}{2} + \frac{c_s}{4}$$

 $d_s :=$ degree of s and $c_s :=$ number of incident corners.

$$\sum_{\mathbf{s}\in\mathcal{S}}\kappa_{\mathbf{s}}=\chi$$

PROOF.
$$\sum_{s \in S} \kappa_s = S - A +$$



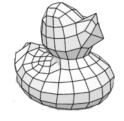
Discrete curvature

$$\kappa_s = 1 - \frac{d_s}{2} + \frac{c_s}{4}$$

 $d_s :=$ degree of s and $c_s :=$ number of incident corners.

$$\sum_{\mathbf{s}\in\mathcal{S}}\kappa_{\mathbf{s}}=\chi$$

PROOF.
$$\sum_{s \in S} \kappa_s = S - A + 4F/4$$



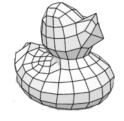
Discrete curvature

$$\kappa_s = 1 - \frac{d_s}{2} + \frac{c_s}{4}$$

 $d_s :=$ degree of s and $c_s :=$ number of incident corners.

$$\sum_{\mathbf{s}\in\mathcal{S}}\kappa_{\mathbf{s}}=\chi$$

PROOF.
$$\sum_{s \in S} \kappa_s = S - A + F$$



Discrete curvature

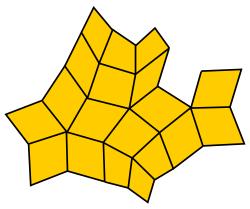
$$\kappa_s = 1 - \frac{d_s}{2} + \frac{c_s}{4}$$

 $d_s :=$ degree of s and $c_s :=$ number of incident corners.

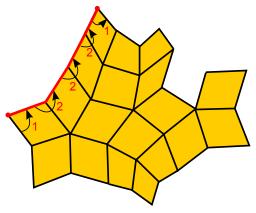
$$\sum_{\mathbf{s}\in\mathcal{S}}\kappa_{\mathbf{s}}=\chi$$

PROOF.
$$\sum_{s \in S} \kappa_s = S - A + F = \chi$$

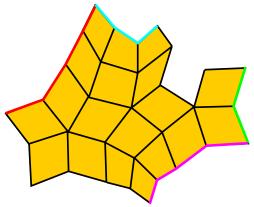
Hyp: All faces are quadrilaterals and all internal vertices have degree ≥ 4 .



Hyp : All faces are quadrilaterals and all internal vertices have degree ≥ 4 .



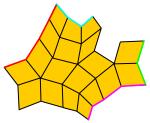
Hyp: All faces are quadrilaterals and all internal vertices have degree ≥ 4 .



4 brackets Theorem (Gersten et Short '90)

The boundary of a non-singular disk has at least 4 brackets.

Hyp: All faces are quadrilaterals and all internal vertices have degree ≥ 4 .

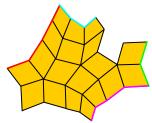


4 brackets Theorem (Gersten and Short '90)

The boundary of a non-singular disk has at least 4 brackets.

PROOF. By Gauss-Bonnet $\sum_{s \in S} \kappa_s = 1$.

Hyp : All faces are quadrilaterals and all internal vertices have degree \geq 4.

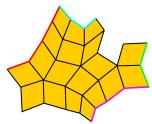


4 brackets Theorem (Gersten and Short '90)

The boundary of a non-singular disk has at least 4 brackets.

PROOF. By Gauss-Bonnet $\sum_{s \in S} \kappa_s = 1$. If s internal: $\kappa_s = 1 - d_s/2 + c_s/4 = 1 - c_s/4 \le 0$.

Hyp : All faces are quadrilaterals and all internal vertices have degree \geq 4.



4 brackets Theorem (Gersten and Short '90)

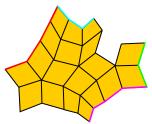
The boundary of a non-singular disk has at least 4 brackets.

PROOF. By Gauss-Bonnet $\sum_{s \in S} \kappa_s = 1$.

If s internal: $\kappa_s = 1 - d_s/2 + c_s/4 = 1 - c_s/4 \le 0$.

If s on the boundary: $\kappa_s = (2 - c_s)/4$.

Hyp: All faces are quadrilaterals and all internal vertices have degree \geq 4.



4 brackets Theorem (Gersten and Short '90)

The boundary of a non-singular disk has at least 4 brackets.

PROOF. By Gauss-Bonnet $\sum_{s \in S} \kappa_s = 1$.

If s internal: $\kappa_s = 1 - d_s/2 + c_s/4 = 1 - c_s/4 \le 0$.

If s on the boundary: $\kappa_s = (2 - c_s)/4$.

Hence, on the boundary: $\#\{s \mid c_s = 1\} \ge \#\{s \mid c_s \ge 3\} + 4$.

Г

Hyp : All faces are quadrilaterals and all internal vertices have degree \geq 4.

Corollary

Every contractible curve (non reduced to a vertex) without spur contains at least 4 brackets.

Hyp : All faces are quadrilaterals and all internal vertices have degree \geq 4.

Corollary

Every contractible curve (non reduced to a vertex) without spur contains at least 4 brackets.

van Kampen '33

Every contractible curve is the label of a reduced disk diagram.

Hyp: All faces are quadrilaterals and all internal vertices have degree ≥ 4 .

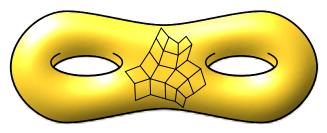
Corollary

Every contractible curve (non reduced to a vertex) without spur contains at least 4 brackets.

van Kampen '33

Every contractible curve is the label of a reduced disk diagram.

PROOF.



Hyp : All faces are quadrilaterals and all internal vertices have degree ≥ 4 .

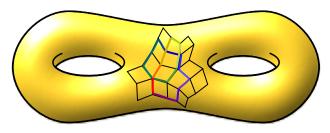
Corollary

Every contractible curve (non reduced to a vertex) without spur contains at least 4 brackets.

van Kampen '33

Every contractible curve is the label of a reduced disk diagram.

PROOF.



Hyp : All faces are quadrilaterals and all internal vertices have degree \geq 4.

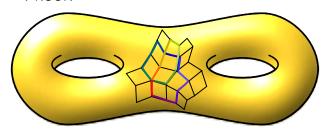
Corollary

Every contractible curve (non reduced to a vertex) without spur contains at least 4 brackets.

van Kampen '33

Every contractible curve is the label of a reduced disk diagram.

PROOF.





Hyp: All faces are quadrilaterals and all internal vertices have degree ≥ 4 .

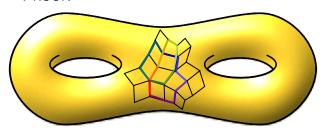
Corollary

Every contractible curve (non reduced to a vertex) without spur contains at least 4 brackets.

van Kampen '33

Every contractible curve is the label of a reduced disk diagram.

PROOF.





Apply the 4 brackets Theorem to this disk.

Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.

The 5 brackets Theorem

The boundary of a non-singular disk with at least one interior vertex has at least 5 brackets.

Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.

The 5 brackets Theorem

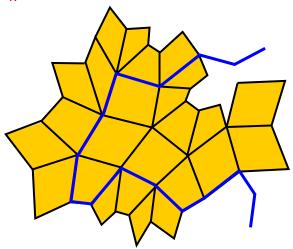
The boundary of a non-singular disk with at least one interior vertex has at least 5 brackets.

PROOF. By Gauss-Bonnet $\sum_{s\in S} \kappa_s = 1$. If s internal: $\kappa_s = 1 - d_s/2 + c_s/4 = 1 - c_s/4 < 0$. If s on the boundary: $\kappa_s = (2 - c_s)/4$.

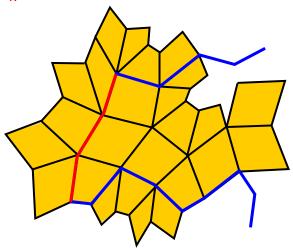
Hence, on the boundary: $\#\{s \mid c_s = 1\} \ge \#\{s \mid c_s \ge 3\} + 5$.

П

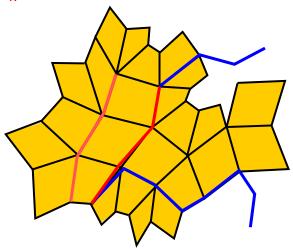
Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.



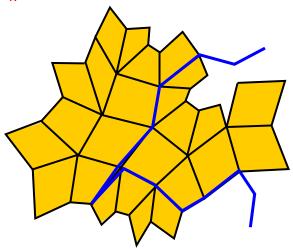
Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.



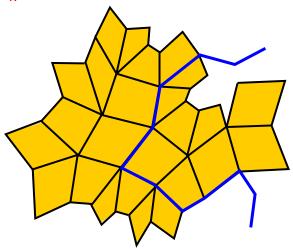
Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.



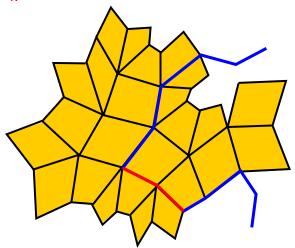
Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.



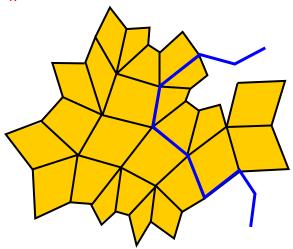
Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.



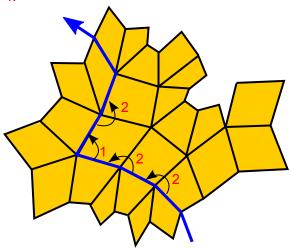
Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.



Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.

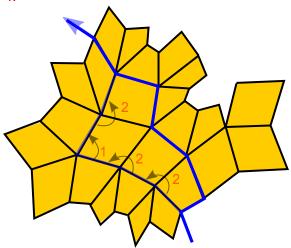


Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.



We shorten by removing brackets and spurs and push to the right.

Hyp: All faces are quadrilaterals and all internal vertices have degree > 4.



We shorten by removing brackets and spurs and push to the right.

L. and Rivaud '12, Erickson and Whittlesey '13

After removing all spurs and brackets and pushing to the right as much as possible, we obtain a canonical representative. It can be computed in linear time.

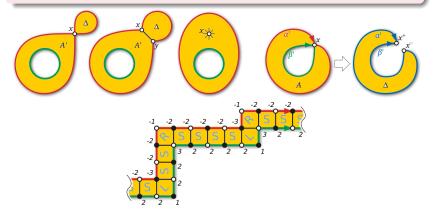
L. and Rivaud '12, Erickson and Whittlesey '13

After removing all spurs and brackets and pushing to the right as much as possible, we obtain a canonical representative. It can be computed in linear time.



L. and Rivaud '12, Erickson and Whittlesey '13

After removing all spurs and brackets and pushing to the right as much as possible, we obtain a canonical representative. It can be computed in linear time.



L. and Rivaud '12, Erickson and Whittlesey '13

After removing all spurs and brackets and pushing to the right as much as possible, we obtain a canonical representative. It can be computed in linear time.

Corollary I

One can decide if two curves are homotopic in linear time.

L. and Rivaud '12, Erickson and Whittlesey '13

After removing all spurs and brackets and pushing to the right as much as possible, we obtain a canonical representative. It can be computed in linear time.

Corollary I

One can decide if two curves are homotopic in linear time.

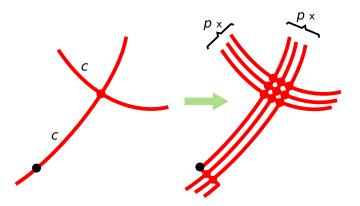
Corollary II

One can compute the primitive root of a curve in linear time.

Non-primitive curves

Lemma

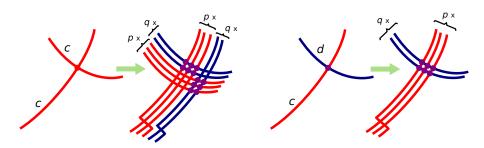
$$i(c^p) = p^2 \times i(c) + p - 1$$



Non-primitive curves

Lemma

$$i(c^p, d^q) = \left\{ egin{array}{ll} 2pq imes i(c) & ext{if } c \sim d ext{ or } c \sim d^{-1}, \\ pq imes i(c, d) & ext{otherwise}. \end{array}
ight.$$



The plan

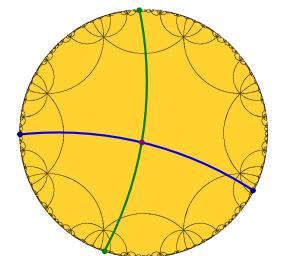
For a given curve c

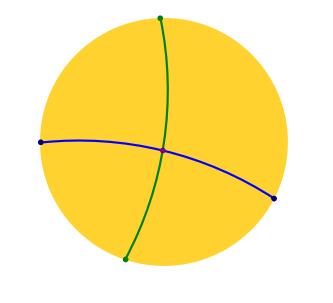
- **1** Determine the primitive root of *c*.
- Count the number of classes of crossing pairs of limit points (for the root of c).
- 3 Use adequate formula if c is not primitive.

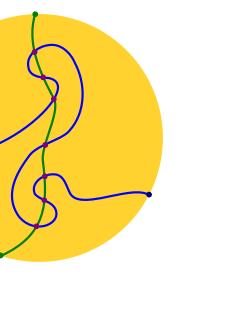
The plan

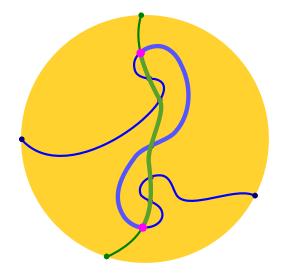
For a given curve c

- Determine the primitive root of *c*.
- Count the number of classes of crossing pairs of limit points (for the root of c).
- Use adequate formula if c is not primitive.



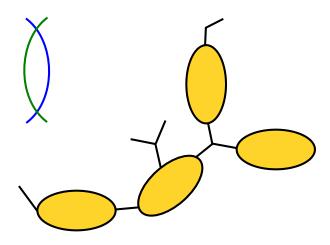


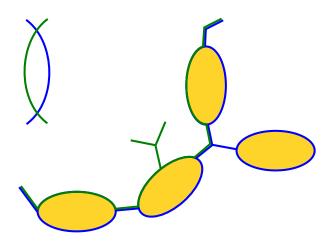


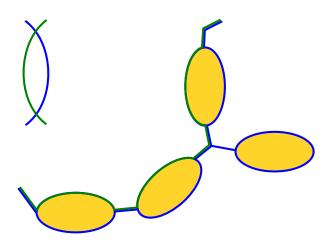


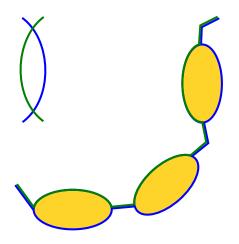
Double paths

A double path is a pair of homotopic paths.

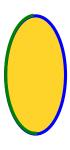


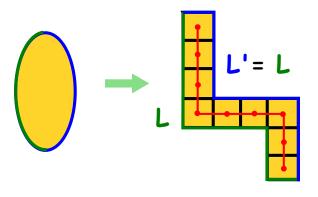


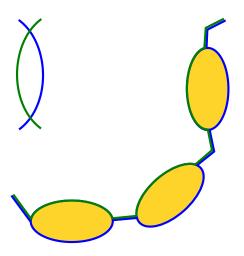


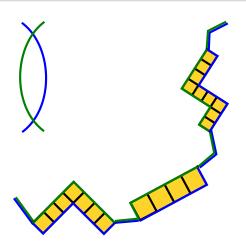


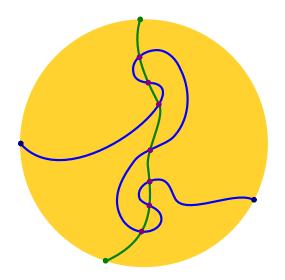


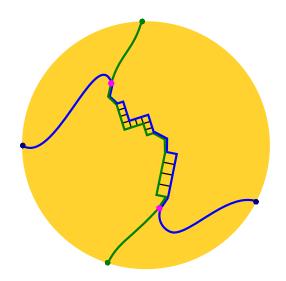












Double paths in a canonical representative (resp. a geodesic) are quasi-flat.

Corollary

Each class of crossing pairs of limit points are uniquely identified by a maximal double path.

Double paths in a canonical representative (resp. a geodesic) are quasi-flat.

Corollary

Each class of crossing pairs of limit points are uniquely identified by a maximal double path.

Lemma

The set of maximal double paths can be computed in quadratic time.

PROOF. A pair of indices (i,j) may occur at most twice in the set of maximal double paths. \Box

Double paths in a canonical representative (resp. a geodesic) are quasi-flat.

Corollary

Each class of crossing pairs of limit points are uniquely identified by a maximal double path.

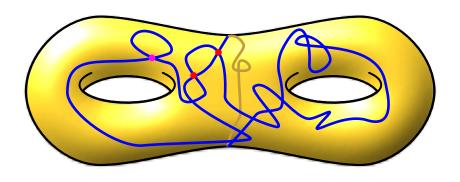
Lemma

The set of maximal double paths can be computed in quadratic time.

Despré and L. '16

The geometric intersection number of one (two) curve(s) can be computed in quadratic time.

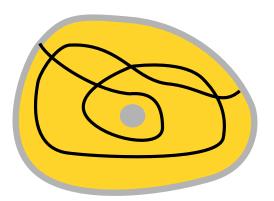
Computing an actual minimal configuration



Hass and Scott '85

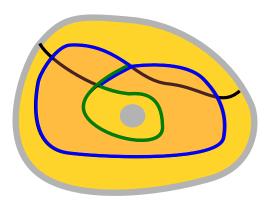
Computing an actual minimal configuration

Hass and Scott '85



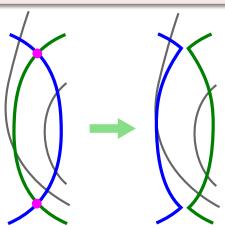
Computing an actual minimal configuration

Hass and Scott '85



Bigon swapping

Hass and Scott '85



Bigon swapping

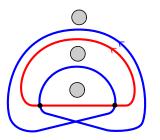
Hass and Scott '85

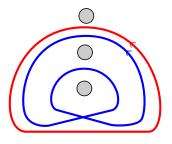
A curve with excess intersection has either a monogon or a singular bigon.

Despré and L. '16

Given c, an homotopic immersion with a minimal number of intersections can be computed in $O(|c|^4)$ time.

Case of two curves





 Propose an algorithm to compute a minimal immersion of two curves.

- Propose an algorithm to compute a minimal immersion of two curves.
- Is quadratic time optimal to just compute the geometric intersection number?

- Propose an algorithm to compute a minimal immersion of two curves.
- Is quadratic time optimal to just compute the geometric intersection number?
- Find a better algorithm (less than quartic) to compute a minimal immersion of single curve.

- Propose an algorithm to compute a minimal immersion of two curves.
- Is quadratic time optimal to just compute the geometric intersection number?
- Find a better algorithm (less than quartic) to compute a minimal immersion of single curve.
- Design a better algorithm to just decide if the geometric intersection number is null.

- Propose an algorithm to compute a minimal immersion of two curves.
- Is quadratic time optimal to just compute the geometric intersection number?
- Find a better algorithm (less than quartic) to compute a minimal immersion of single curve.
- Design a better algorithm to just decide if the geometric intersection number is null.
- Propose an algorithm to compute a minimal immersion in hyperbolic configuration (cf. Hass and Scott '99).

Thank you for your celtantion