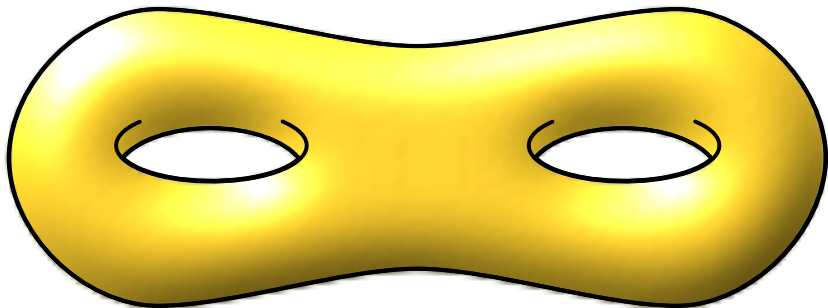


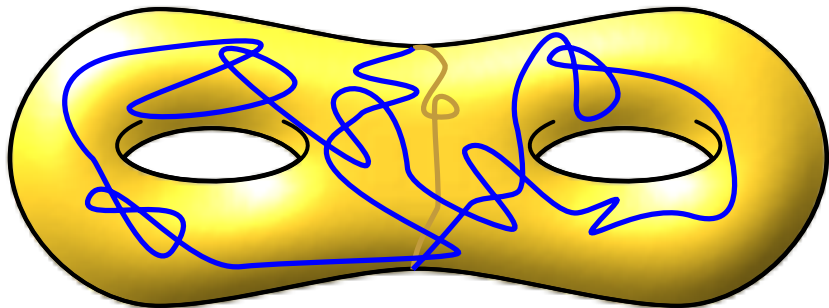
Autour du nombre géométrique d'intersection

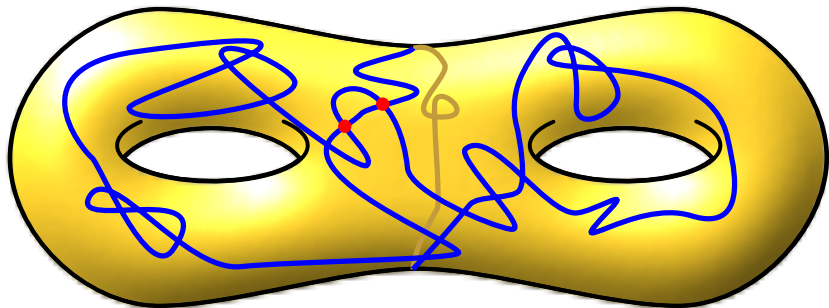
Francis Lazarus
GIPSA-Lab, Grenoble, CNRS

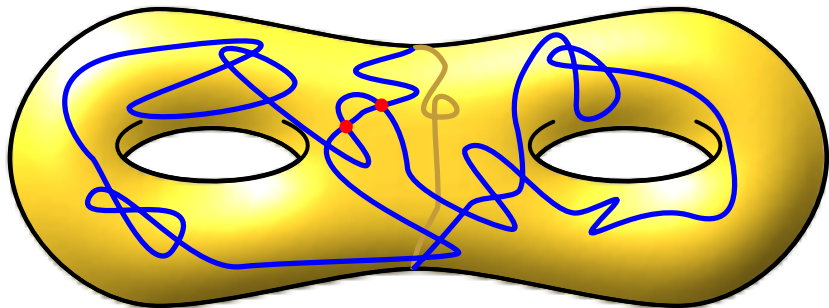


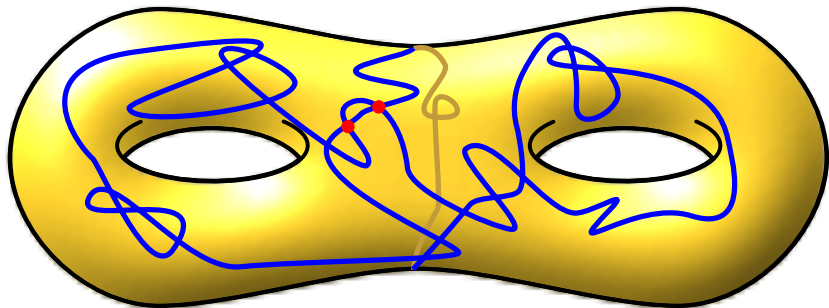
Lionel Walden, Les Docks de Cardiff, 1894

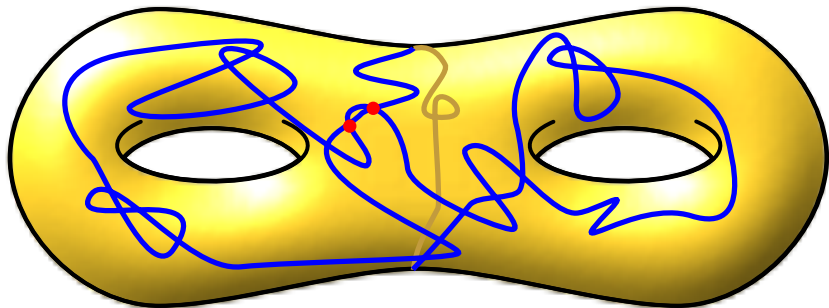


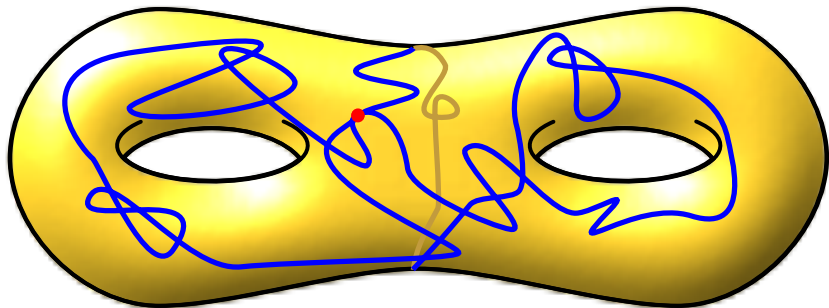


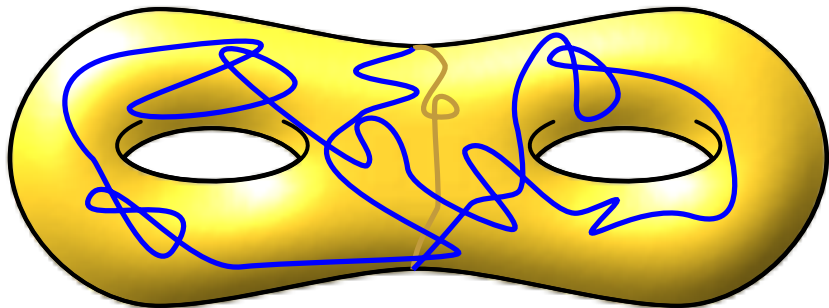


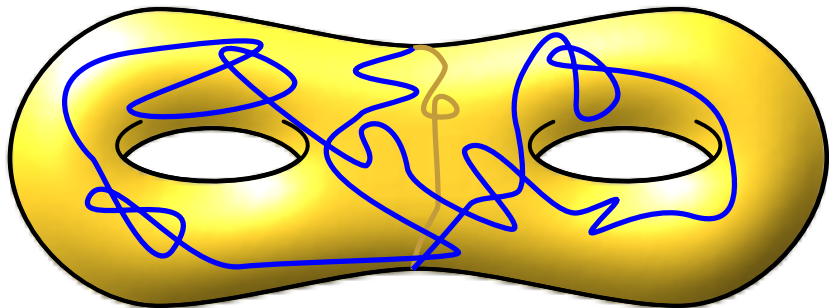


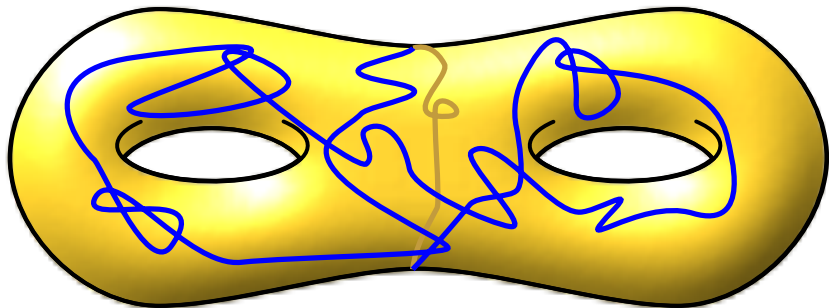


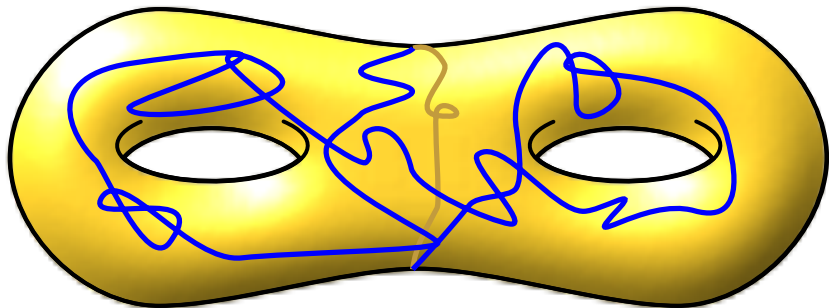


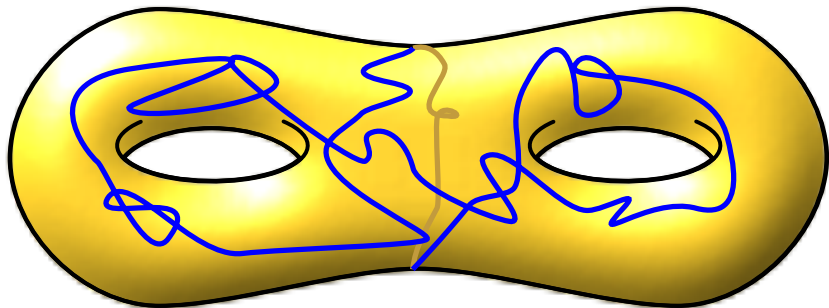


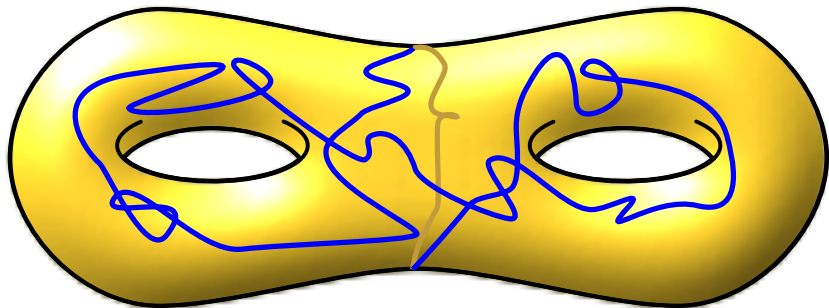


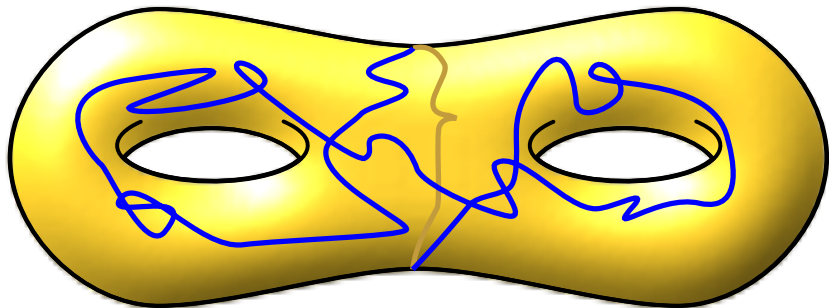


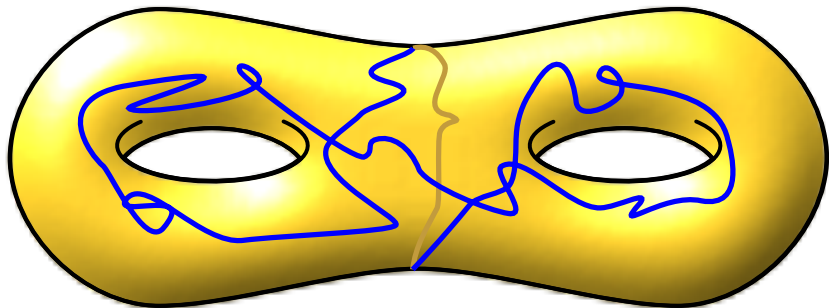


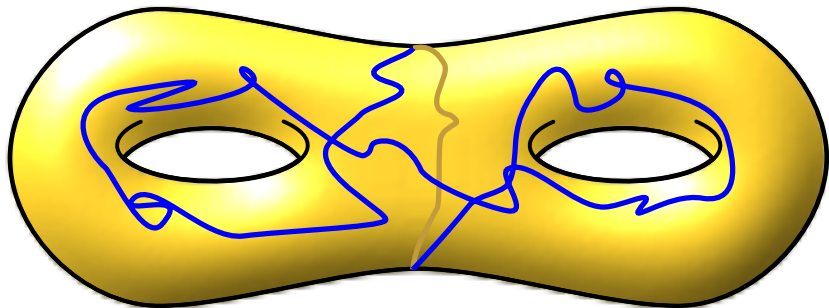


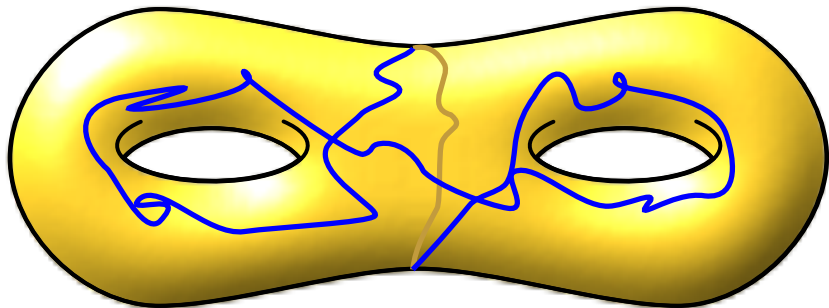


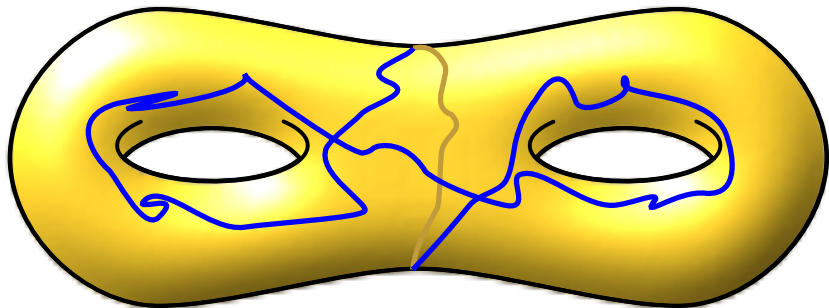


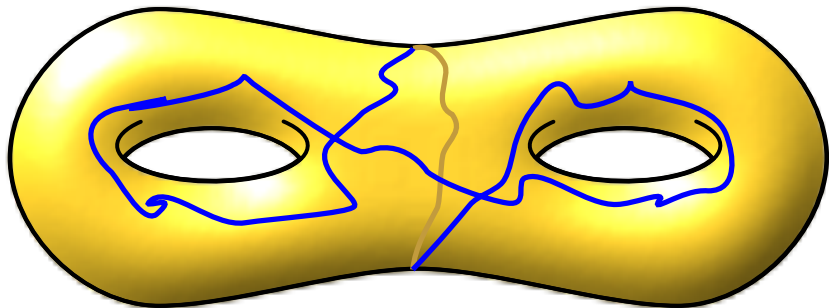


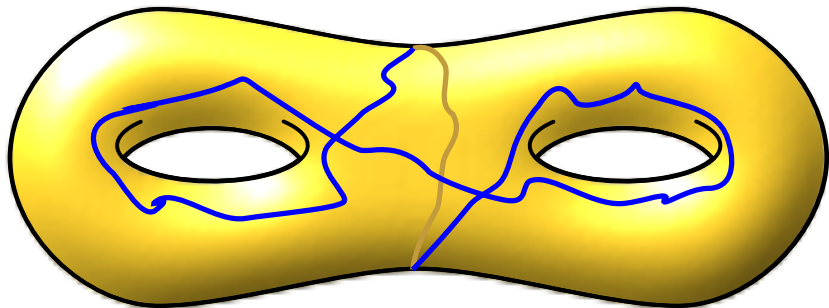


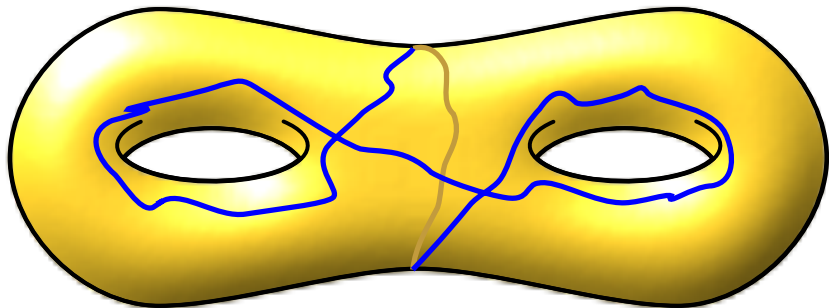


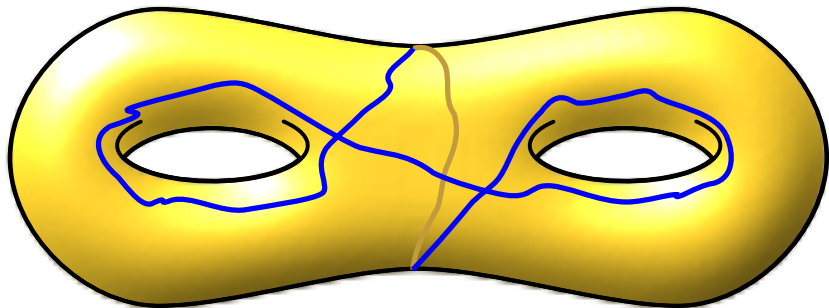


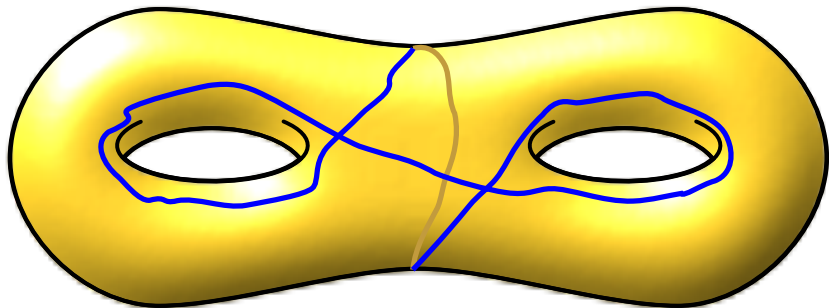


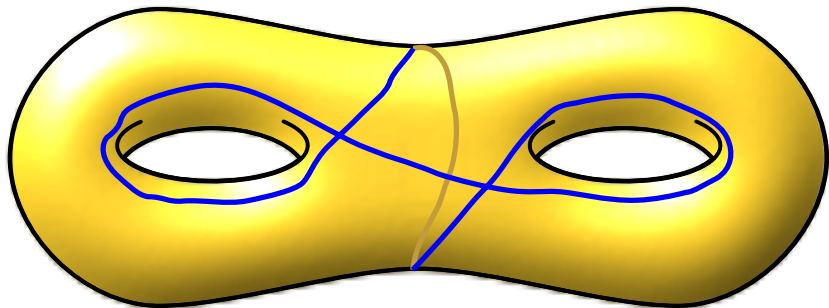


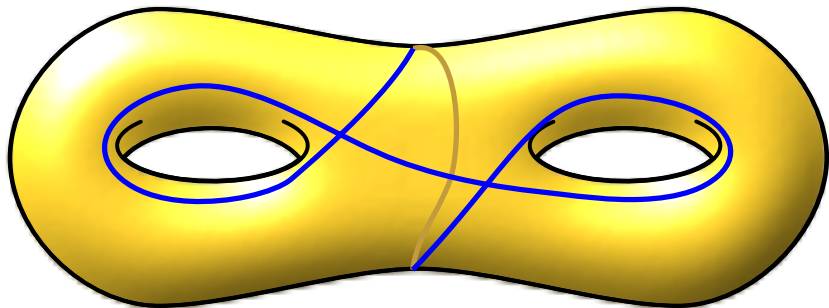


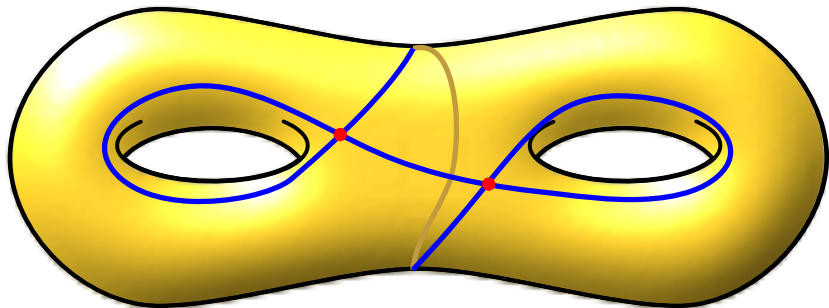


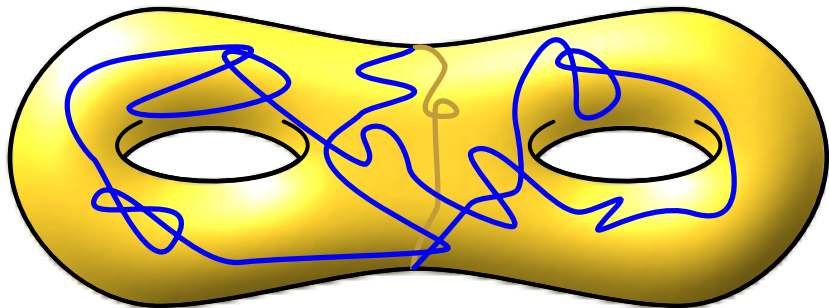


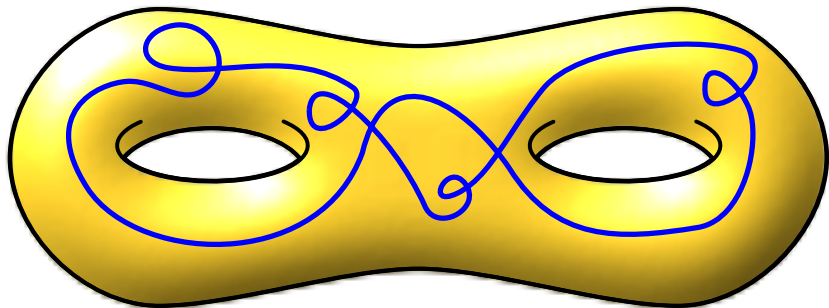


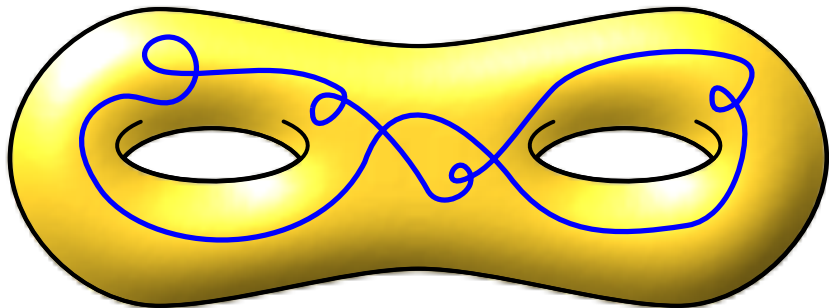


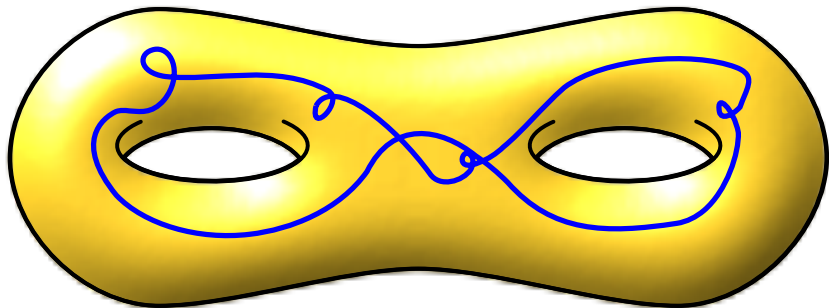


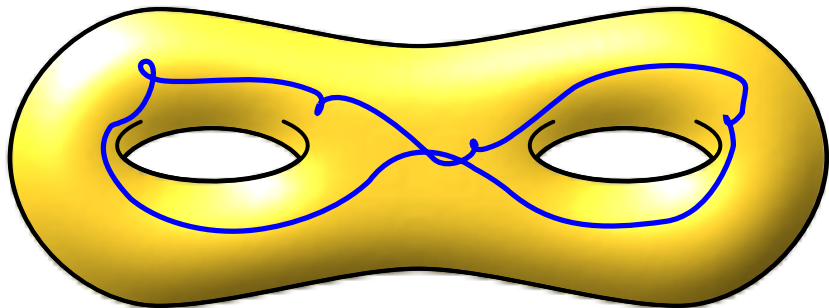


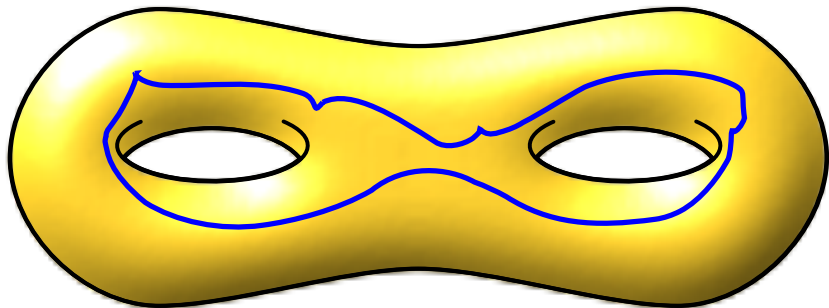


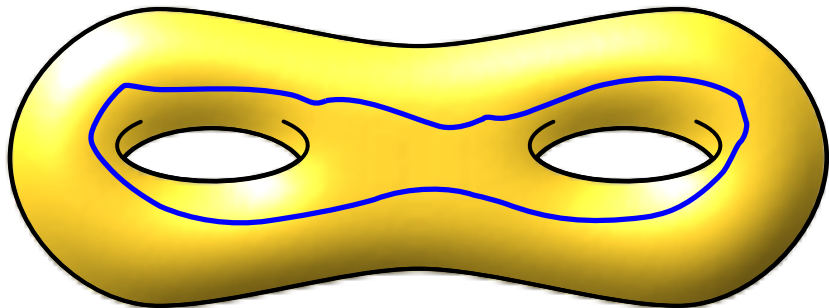


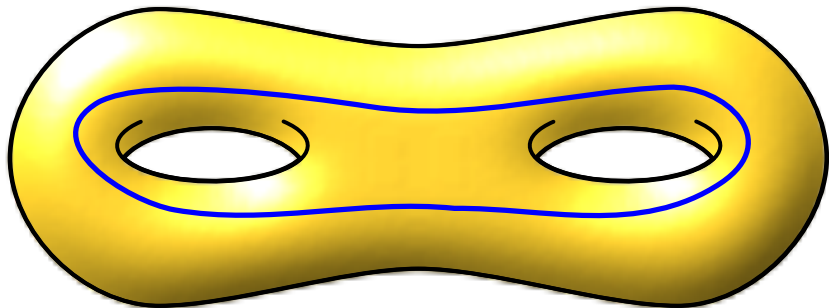


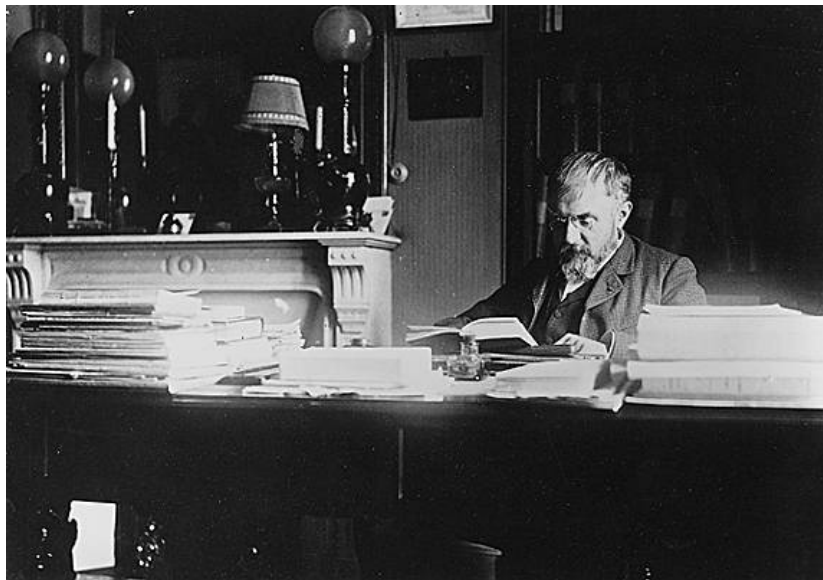












1854 – 1912

SECOND COMPLÉMENT

Proceedings of the London Mathematical Society, t. IV, n. 303-324 (15 juin 1900).

1. — Introduction

Dans le *Journal de l'École Polytechnique* (volume du centenaire de la fondation de l'École, 1894) j'ai publié en Mémorial *Analystis nris*, où j'étudie les variétés de l'espace à plus de trois dimensions et les propriétés des nombres de Betti. C'est à ce Mémorial que se rapportent les renvois que j'ai cités ci-dessus à faire fréquemment dans la suite, en mentionnant seulement le titre *Analystis nris*.

Le même théorème a été énoncé par M. Picard dans sa *Théorie des fonctions*

M. Hoegedahl vient de revenir sur ce même problème dans un travail très remarquable, publié en langue danoise, sous le titre « *Forestudier til en topologisk teori for de algebraiske Sammenhæng* » (Copenhague, des Nordiske Forlag Ernst Bejstien, 1898). D'après lui le théorème en question est inexact et les démonstrations sont sans valeur.

Avant d'examiner les objections de M. Hensgaard, il convient de faire une distinction. Il y a deux manières de définir les nombres de Betti.

qu'en ne puisse pas trouver de variété à $p+1$ dimensions faisant partie de V et dont v_1, v_2, \dots, v_n constituent le feuilletre complète; mais que, si on leur adjoint une $(n+1)^{\text{ème}}$ variété à p dimensions que j'appellerai v_{n+1} , et qui for-

8178

LES CYCLES DES SURFACES ALGÈBRIQUES:

QUATRIÈME COMPLÉMENT A L'ANALYSIS SITUS

L'ANALYSIS SITUS

L'ANALYSIS SITUS

L'ANALYSIS SITUS

Journal de Mathématiques, t. 8, p. 189-212 (1901)

8.1. Introduction

Les beaux travaux de M. Pinard sur les Surfaces algébriques ont mis depuis longtemps en évidence l'importance de la notion des cycles à une, deux et trois dimensions. Tel ou tel cycle apparaît naturellement à une certaine

principes que j'ai exposés dans l'Analyse situs et ses deux premiers compléments (*Journal de l'École Polytechnique*, tome du centenaire; *Rendiconti del Circolo Matematico di Palermo*, t. XIII; *Proceedings of the London*

Mathematical Society, vol. XXXII) et j'ai obtenu ainsi certains résultats partiels que j'ai déjà annoncés dans une Note aux Comptes rendus et qui méritent cependant une nouvelle mise au point. M. Borel.

Je rappelle qu'étant donné une variété V fermée à p dimensions, je trace sur cette variété d'autres variétés, fermées ou non, d'un moins grand nombre

Si ΣW_q est un ensemble de variétés à q dimensions et ΣW_{q-1} un ensemble

de variétés à $g-1$ dimensiones, la congruence

$$EW_g \equiv EW_{g-1}$$

signifie (par définition) que $\mathbf{I}W_{q-1}$ forme la frontière complète de l'ensemble

Nous sommes donc conduit à une contradiction, ce qui veut dire que l'hypothèse faite au début était absurde et que *le cycle K doit être bouclé*.

C. Q. F. D.

4. Nous avons vu au paragraphe précédent qu'il est relativement aisé de reconnaître si un cycle donné est homologue à un cycle non bouclé, ou si deux cycles donnés sont respectivement homologues à deux cycles qui ne se coupent pas. Nous allons dans le présent paragraphe examiner une question analogue :

Comment reconnaître si un cycle donné est équivalent à un cycle non bouclé, ou si deux cycles donnés sont équivalents à deux cycles qui ne se coupent pas ?

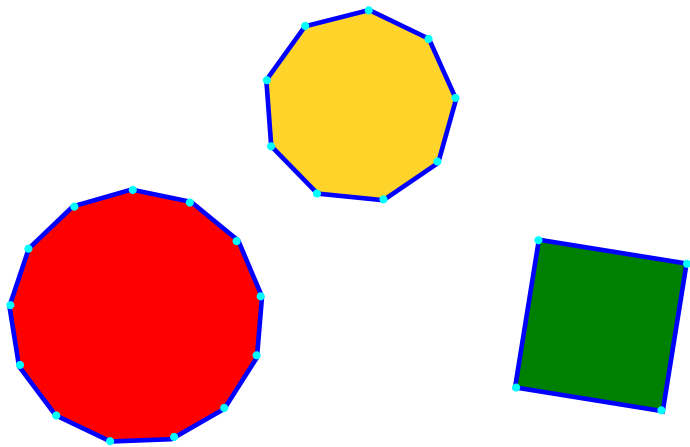
Mais avant d'aborder cette question, revenons sur la définition de l'équivalence.

Jusqu'ici nous avons toujours entendu cette équivalence de la façon suivante :

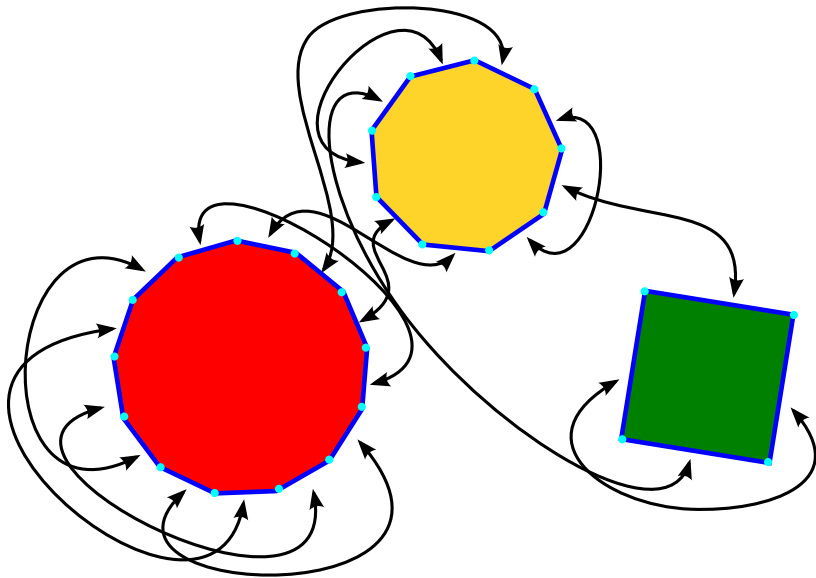
Quand nous écrivons

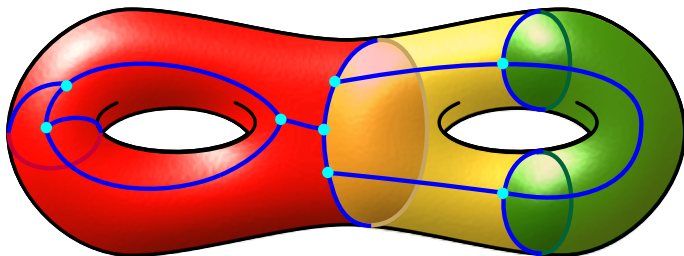
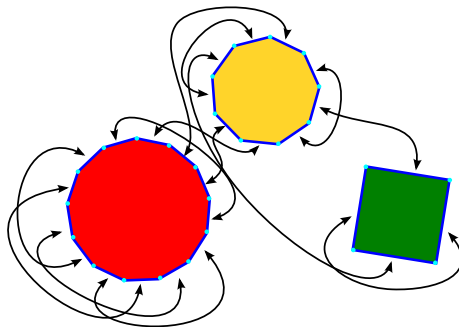
$$C \equiv C',$$

What is a surface?



What is a surface?





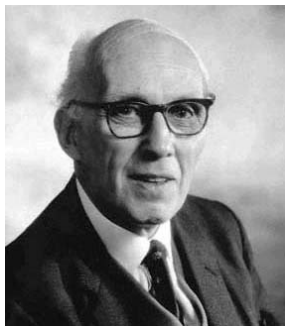
Über den Begriff der Riemannschen Fläche.

Von TIBOR RADÓ in Szeged.

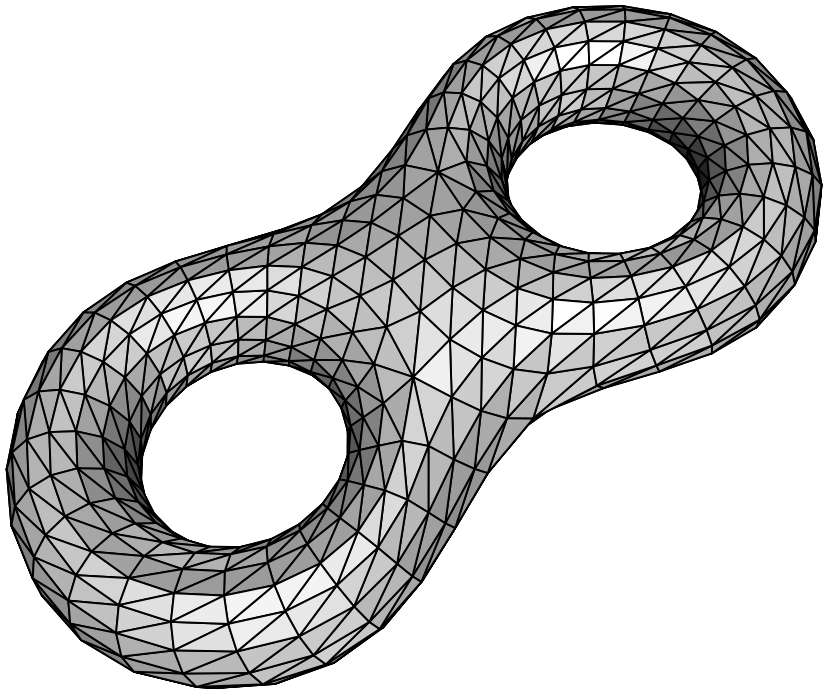
Einleitung.

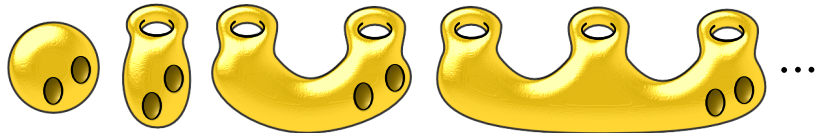
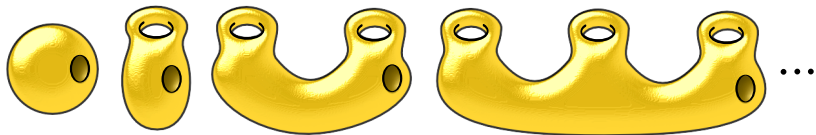
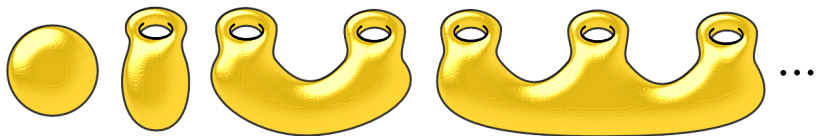
Die vorliegende Arbeit enthält eine Untersuchung, zu welcher ich beim Studium des grundlegenden Werkes des Herrn WEYL über *Die Idee der Riemannschen Fläche* geführt wurde. Bekanntlich wird in diesem Buche der Begriff der RIEMANNschen Fläche zum ersten Male in vollkommen strenger Weise erklärt, und zwar wie folgt. Eine RIEMANNsche Fläche ist eine zweidimensionale Mannigfaltigkeit, welche trianguliert werden kann, und für welche eine konforme Abbildung im Kleinen mitgegeben ist. Wir werden diese Begriffe eingehend besprechen (§§ 1 und 4), müssen aber gleich an dieser Stelle die Forderung der *Triangulierbarkeit* doch etwas genauer betrachten, um unser Problem formulieren zu können.

Der Ausdruck *zweidimensionale Mannigfaltigkeit* wird nicht von allen Autoren in demselben Sinne gebraucht. Sie ist jedenfalls ein zusammenhängender topologischer Raum im HAUSDORFFschen Sinne, welcher im Kleinen der xy -Ebene homöomorph ist; aber es wird manchmal noch die Forderung an sie gestellt, sie soll dem zweiten HAUSDORFFschen Abzählbarkeitsaxiom genügen. Wollte man den Ausdruck zweidimensionale Mannigfaltigkeit bei der Erklärung der RIEMANNschen Fläche in diesem Sinne verstehen, so wäre die Forderung der Triangulierbarkeit überflüssig. Unter Voraussetzung dieses Abzählbarkeitsaxioms bietet nämlich die Triangulierung einer zweidimensionalen Mannigfaltigkeit keine prinzipielle, sondern nur technische Schwierigkeiten, und die explizite Forderung der Triangulierbarkeit würde einfach bedeuten, dass man mit möglicherweise umständlichen, aber im Grunde ganz einfachen Betrachtungen keine Zeit verlieren will.



Tibor Radó
1895 – 1965



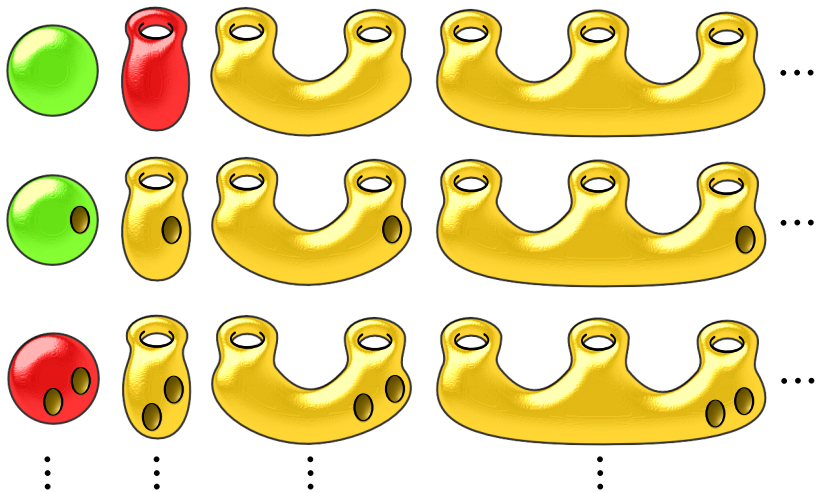


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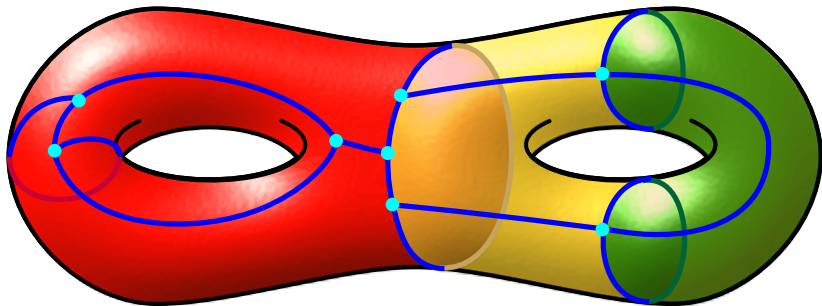
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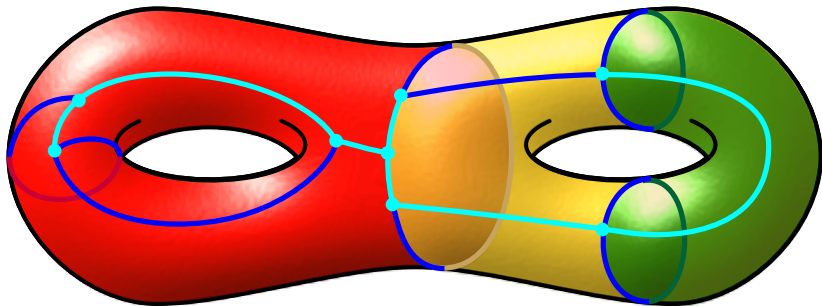
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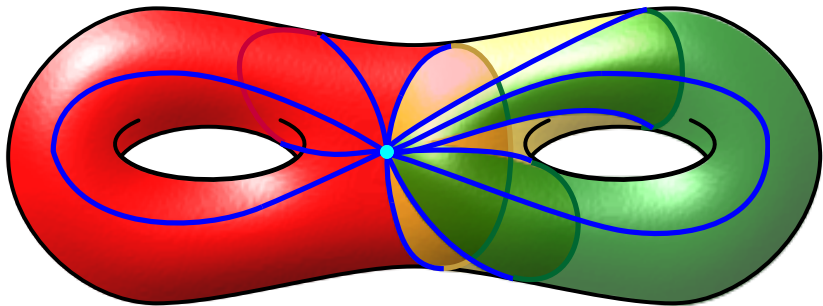
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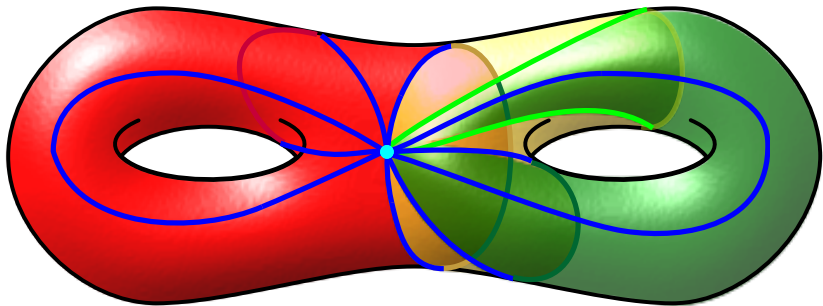


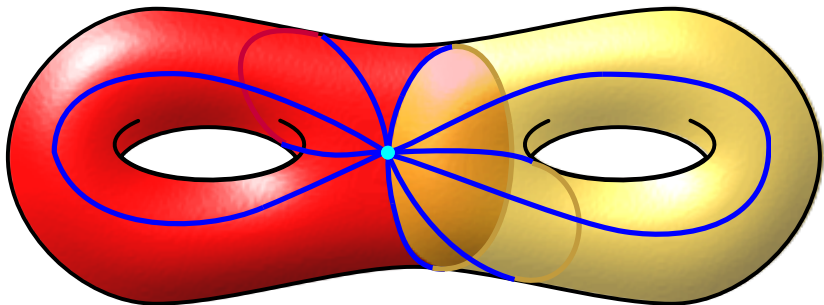
$$S - A + F = \chi = 2 - 2g - b$$

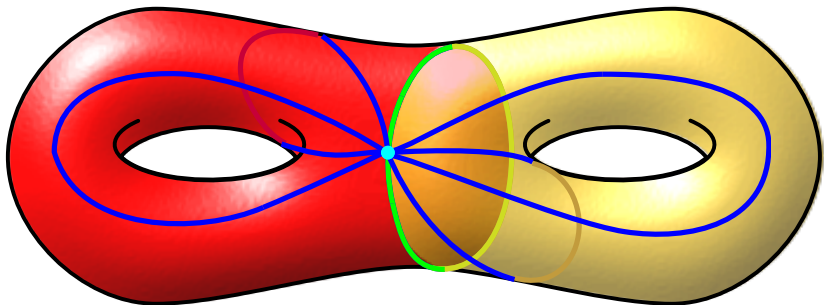


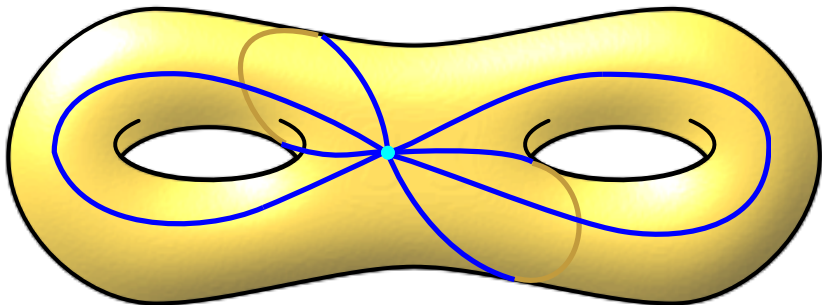


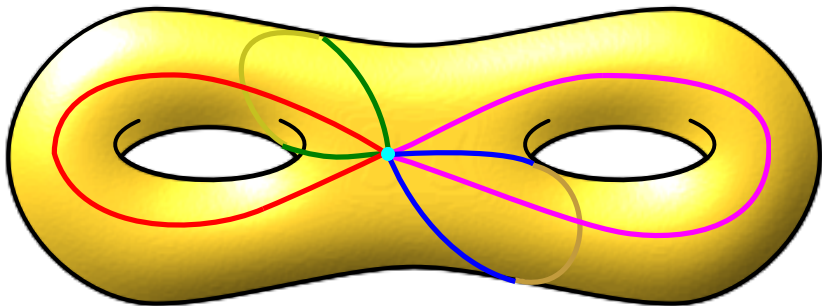


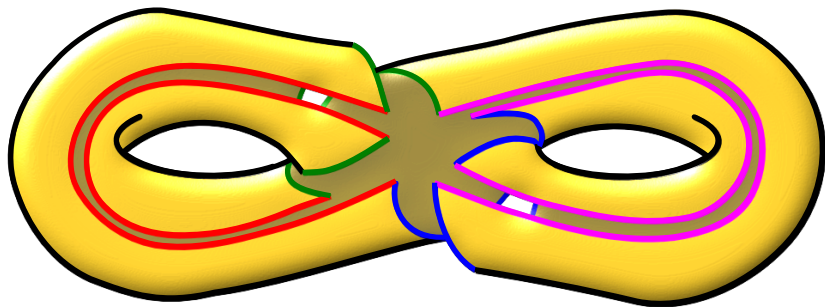


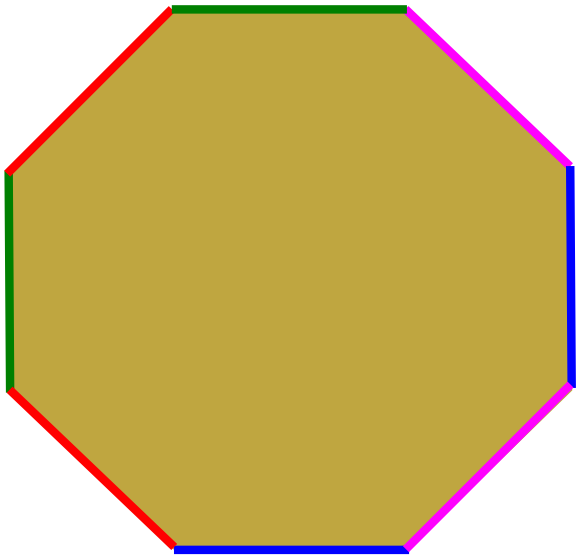




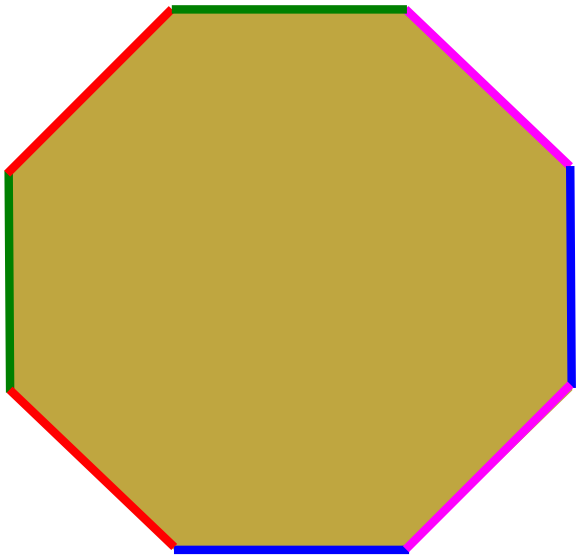


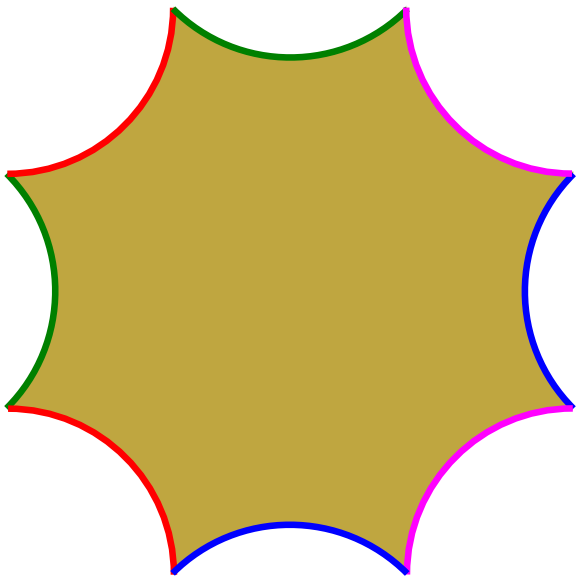






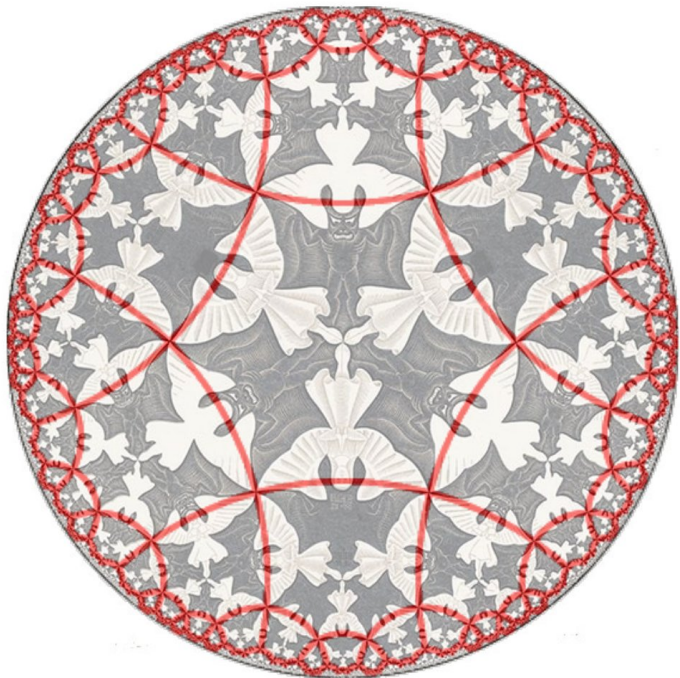
On peut se placer, dans l'étude de la question qui nous occupe, à plusieurs points de vue différents. Représentons d'abord notre surface par un polygone fuchsien R_0 de la première famille, construisons les différents transformés de ce polygone par les transformations du groupe fuchsien correspondant G ; ces transformés rempliront le cercle fondamental. Un cycle quelconque C sera alors représenté par un arc de courbe MM' , allant d'un point M à un de ses transformés M' . Deux cycles proprement équivalents seront représentés par deux arcs de courbe MPM' et MQM' ayant mêmes extrémités et réciproque-

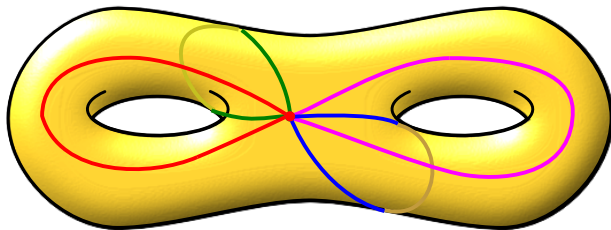
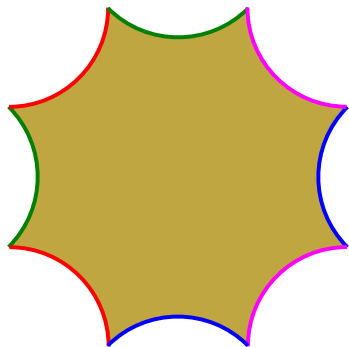


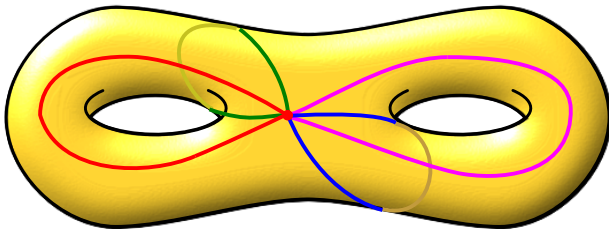
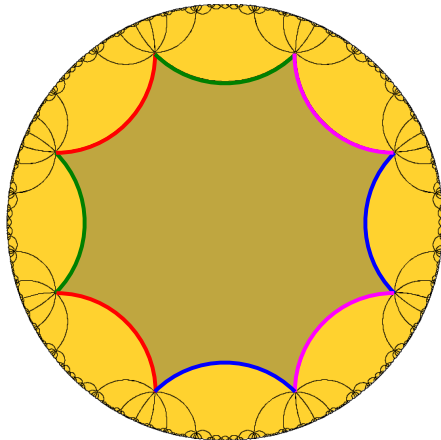


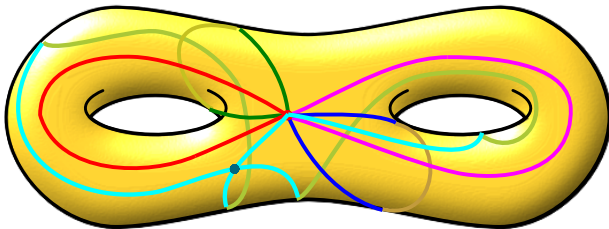
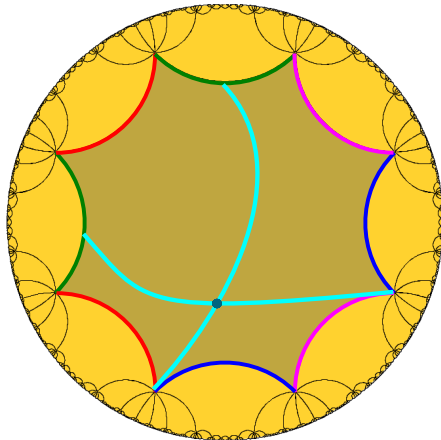


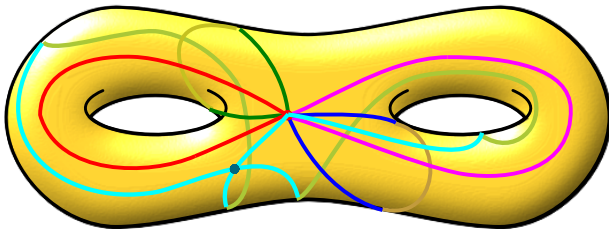
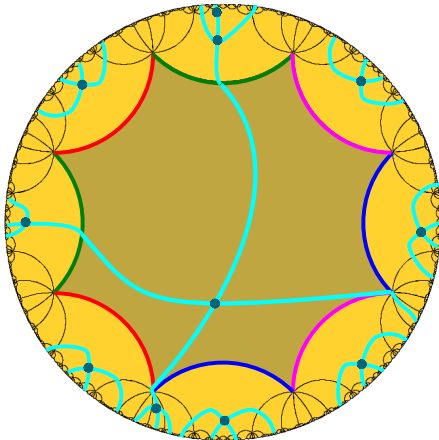


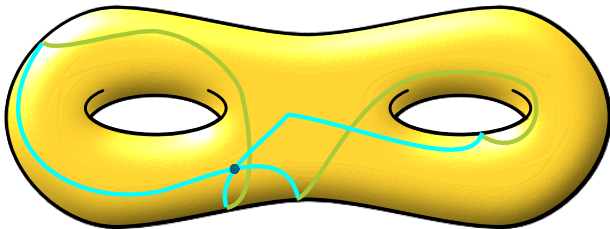
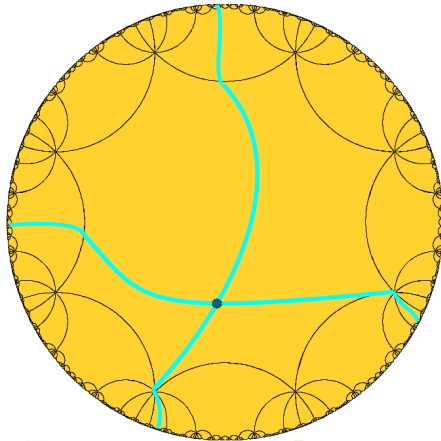


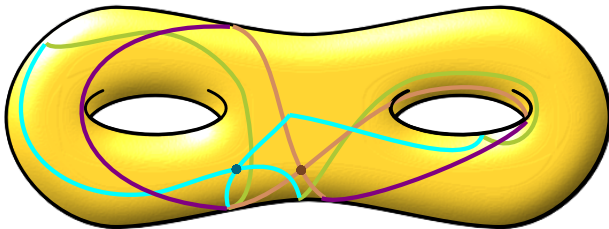
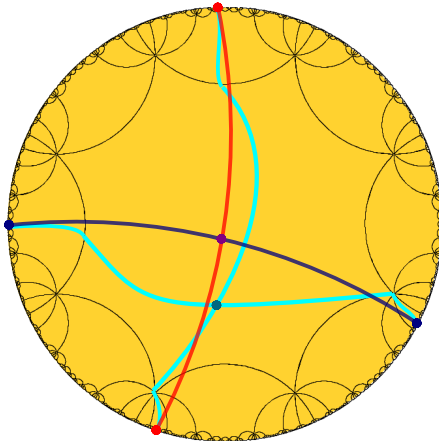


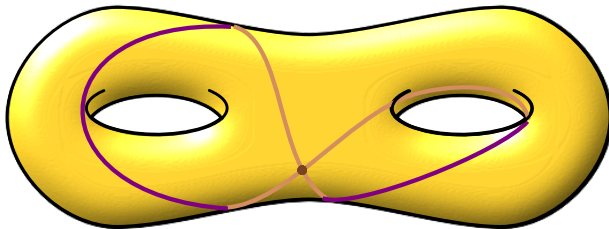
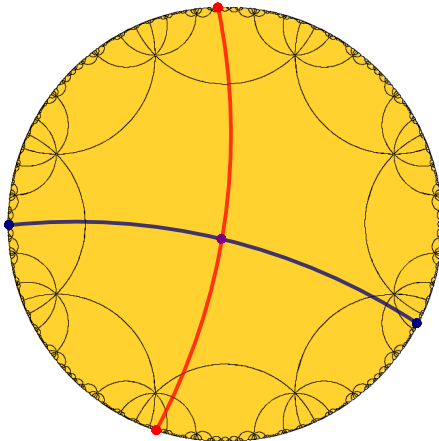












ALGORITHMS FOR JORDAN CURVES ON COMPACT SURFACES

BY BRUCE L. REINHART*

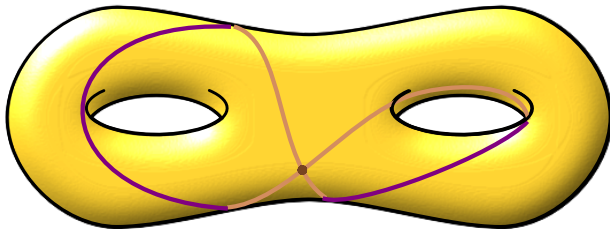
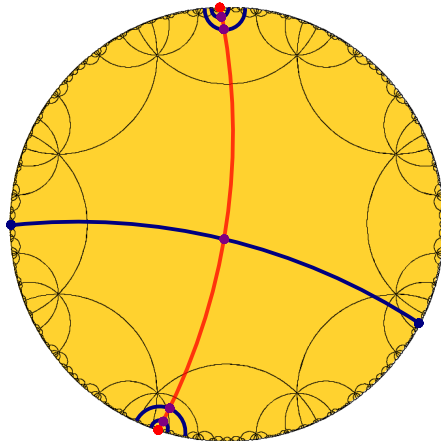
(Received October 20, 1960)

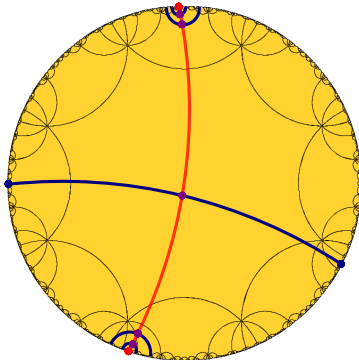
As is well known, free homotopy classes of mappings of the circle into a space correspond to conjugacy classes in the fundamental group. If the space is two-dimensional, such a free homotopy class may or may not contain a simple closed (that is, Jordan) curve. Our purpose here is to give an algorithm for determining which free homotopy classes admit such a curve in the case of compact surfaces of negative Euler number. Our algorithm is applicable to orientable and non-orientable surfaces, with or without boundary. This problem was first discussed by Dehn [4] and later by Baer [1] and Goeritz [5]. All of these studies are based on cutting the surface into spheres with 3 holes and deriving a (hopefully) canonical form for each free homotopy class by pasting together curves on each sphere with holes. Except for the surface of genus 2 without boundary, no complete answer has been achieved in this way. We shall treat the problem globally by imbedding the fundamental group into the group of motions of the hyperbolic plane, in the spirit of Poincaré [14] and numerous works of Nielsen. As a preliminary, we give in §1 a simple method for assigning to each word in the usual generators of the fundamental group a curve on the surface which has double points only in the neighborhood of the base point, and no other multiple points. We call such a curve an indicating curve for the word. In §2, we apply the fact that each motion of the hyperbolic plane which arises in our problem leaves fixed a unique geodesic, its axis. The projection of an axis onto the surface is the unique geodesic lying in the free homotopy class containing the motion of which it is axis. For orientable surfaces, a free homotopy class admits a simple closed curve if and only if its geodesic is simple [14, p. 467], while

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	closed surface	counting	free homotopy	special feature
Chillingworth '69				winding number
Birman & Series '84			✓	retraction onto a graph
Cohen & Lustig '87		✓	✓	retraction onto a graph
Lustig '87	✓	✓	✓	canonical representative
de Graaf & Schrijver '97	✓	✓	✓	Reidemeister moves
Paterson '02	✓	✓	✓	Reidemeister moves
Gonçalves et al. '05	✓	✓	✓	algebraic approach





- If a curve c is **primitive** its lifts are uniquely defined by their limit points.
- If τ is the hyperbolic translation corresponding to a lift \tilde{c}_0 of a primitive c then

$$i(c) = |\{\text{set of pairs of limit points crossing } \tilde{c}_0\} / \langle \tau \rangle|$$

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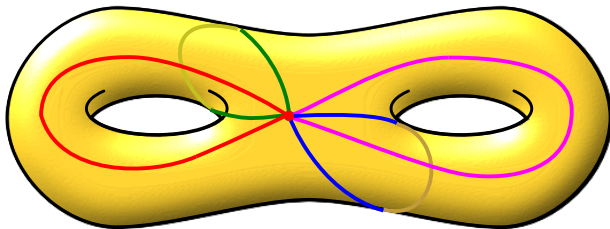
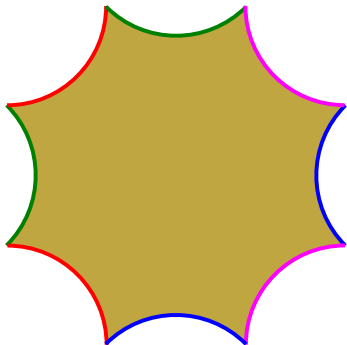
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The plan

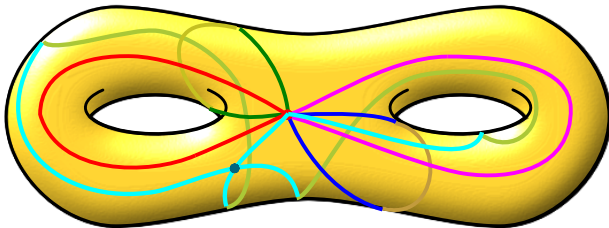
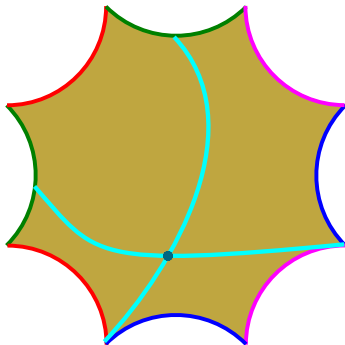
For a given curve c

- 1 Determine the primitive root of c .
- 2 Count the number of classes of crossing pairs of limit points (for the root of c).
- 3 Use adequate formula if c is not primitive.

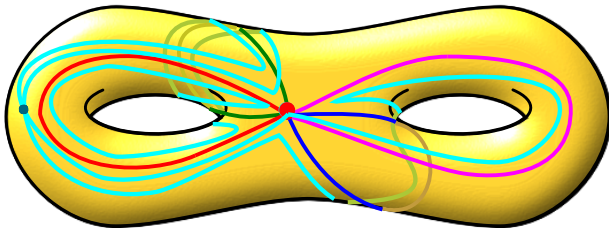
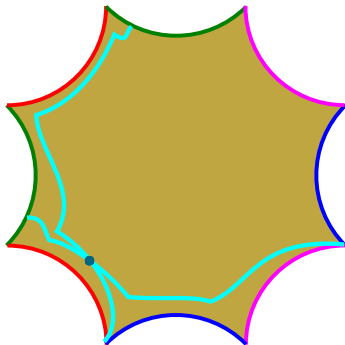
A combinatorial framework



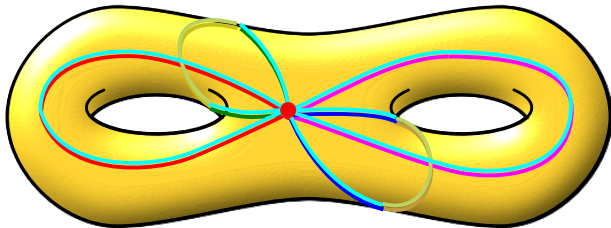
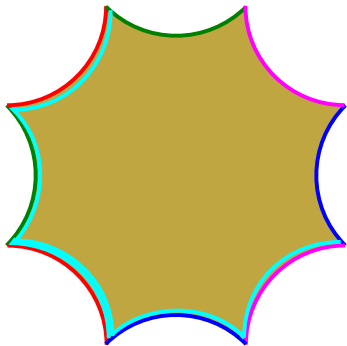
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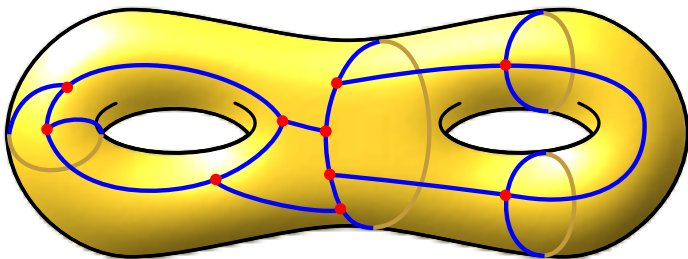
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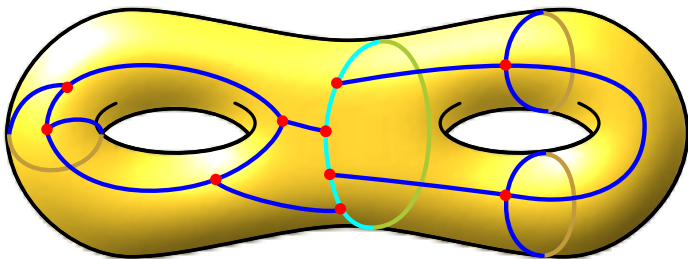
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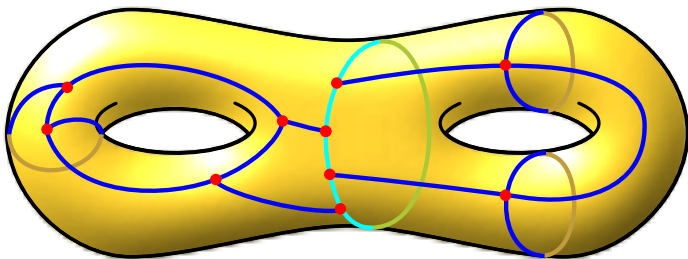
A combinatorial framework: elementary homotopies



A combinatorial framework: elementary homotopies

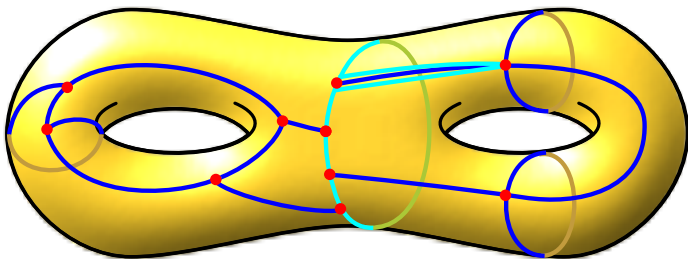


A combinatorial framework: elementary homotopies



Two rules:

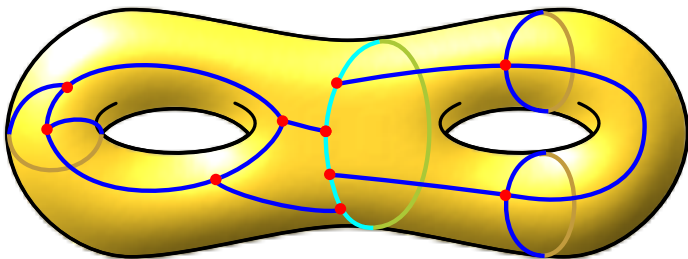
A combinatorial framework: elementary homotopies



Two rules:

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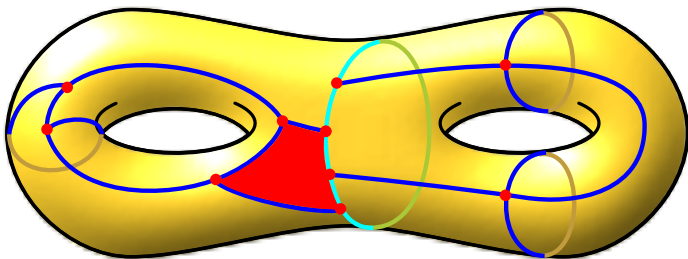
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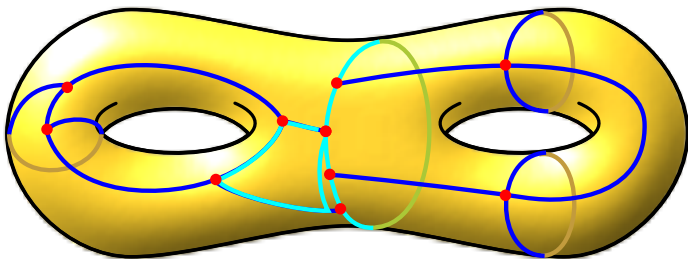
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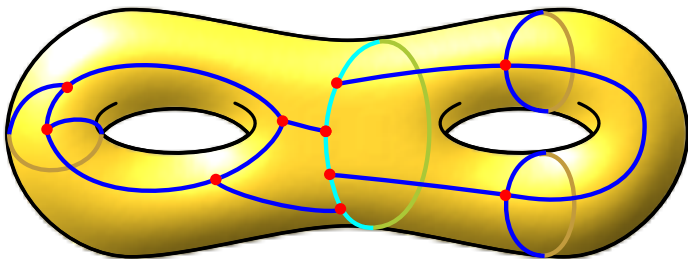
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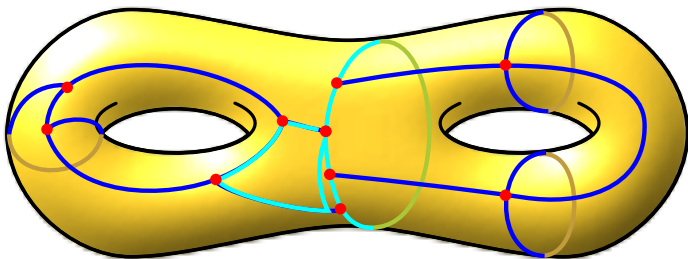
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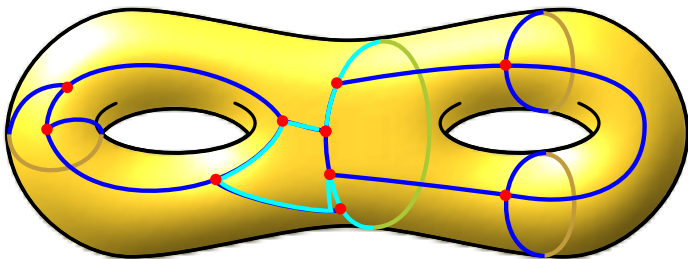
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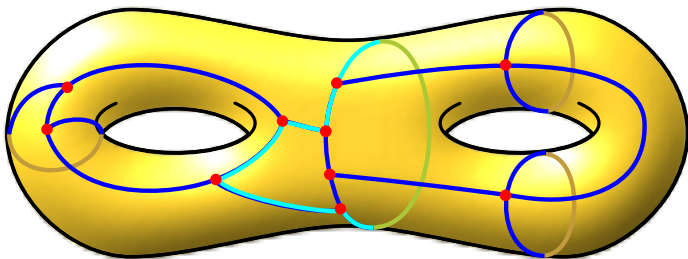
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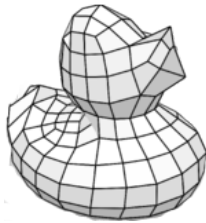
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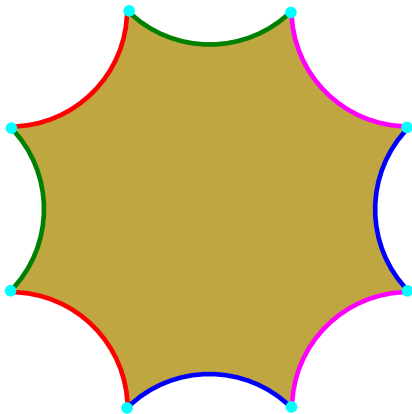
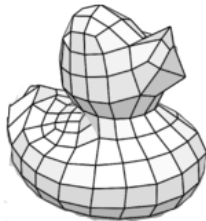
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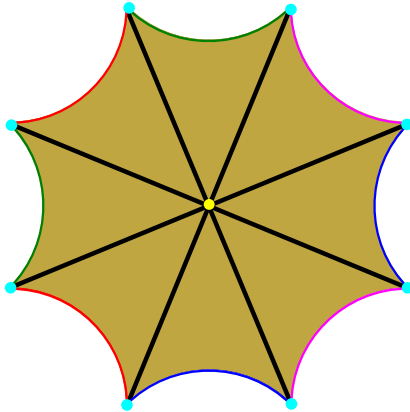
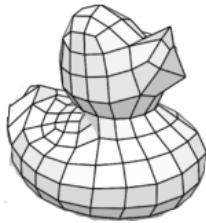
Hyp : All faces are quadrilaterals.



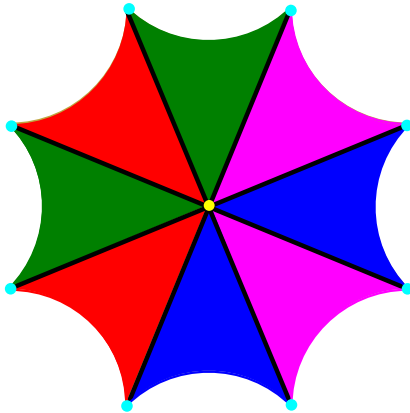
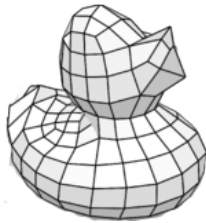
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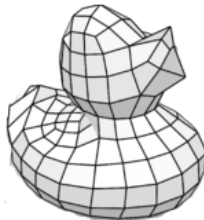
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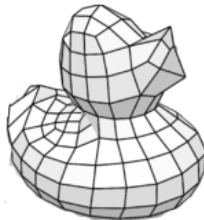


Discrete curvature

$$\kappa_s = 1 - \frac{d_s}{2} + \frac{c_s}{4}$$

d_s := degree of s and c_s := number of incident corners.

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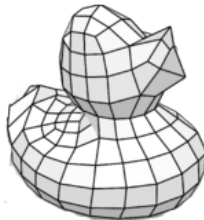
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Combinatorial Gauss-Bonnet Theorem

$$\sum_{s \in S} \kappa_s = \chi$$

PROOF. $\sum_{s \in S} \kappa_s =$

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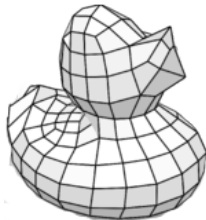
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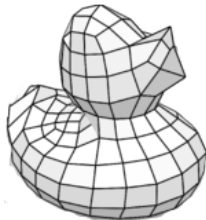
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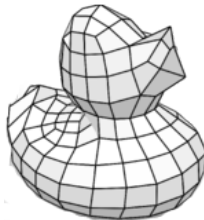
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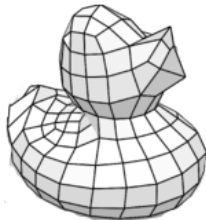
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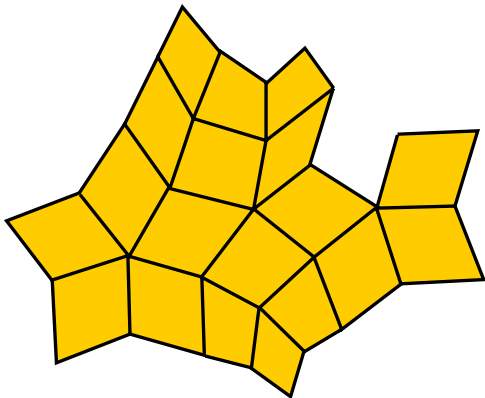
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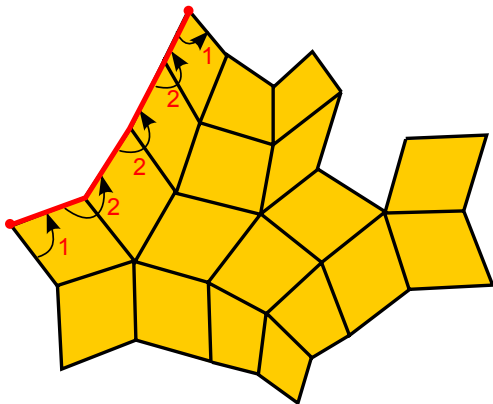
The 4 brackets Theorem

Hyp : All faces are quadrilaterals and all internal vertices have degree ≥ 4 .



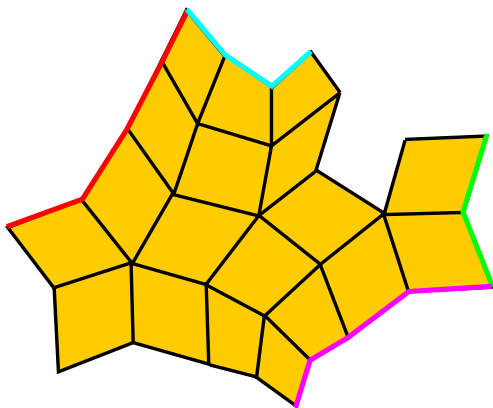
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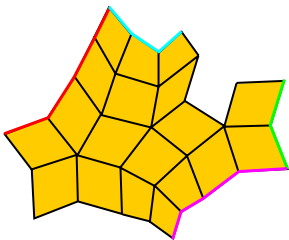


4 brackets Theorem (Gersten et Short '90)

The boundary of a non-singular disk has at least 4 brackets.

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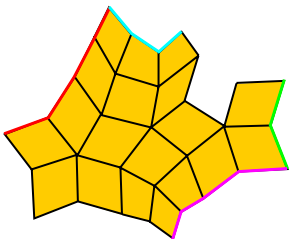
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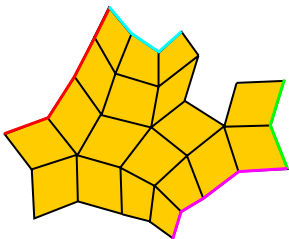
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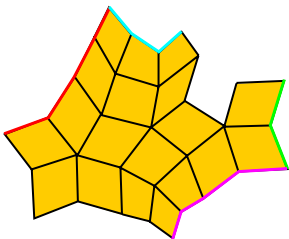
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Hence, on the boundary: $\#\{s \mid c_s = 1\} \geq \#\{s \mid c_s \geq 3\} + 4$.



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Corollary

Every contractible curve (non reduced to a vertex) without spur contains at least 4 brackets.

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van Kampen '33

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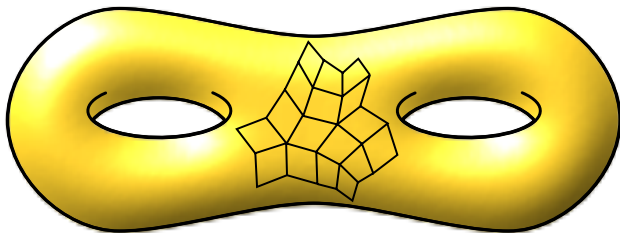
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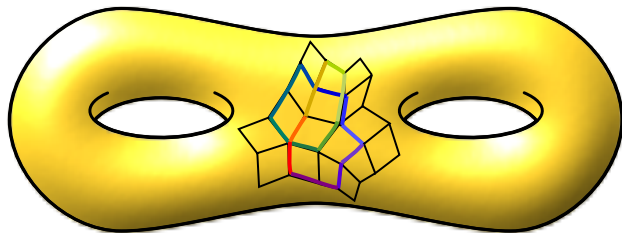
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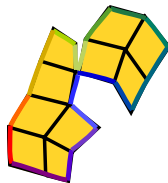
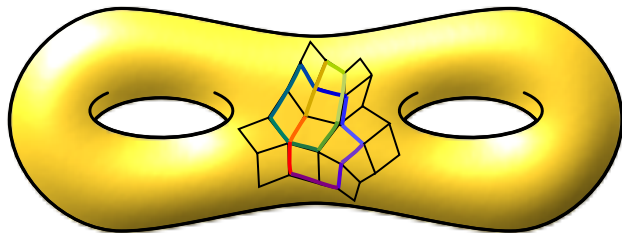
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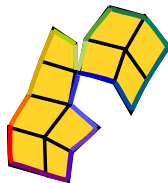
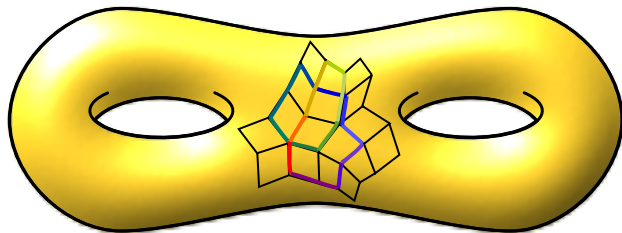
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Apply the 4 brackets Theorem to this disk. \square

Hyp : All faces are quadrilaterals and all internal vertices have degree > 4 .

The 5 brackets Theorem

The boundary of a non-singular disk with at least one interior vertex has at least 5 brackets.

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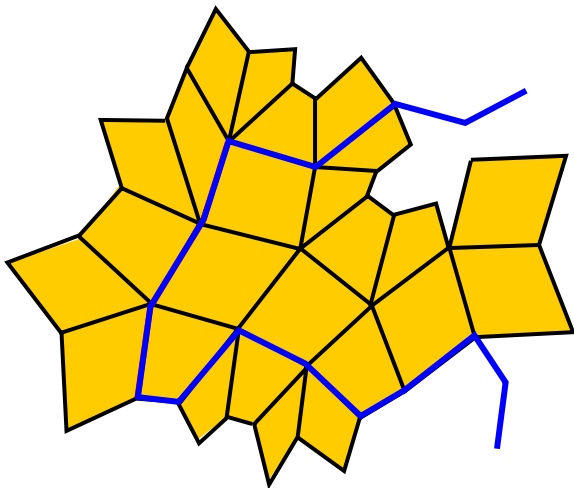
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□

Canonical representative

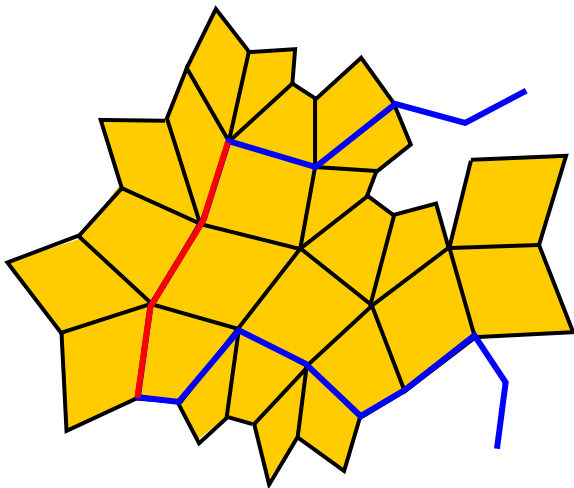
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We shorten by removing brackets and spurs

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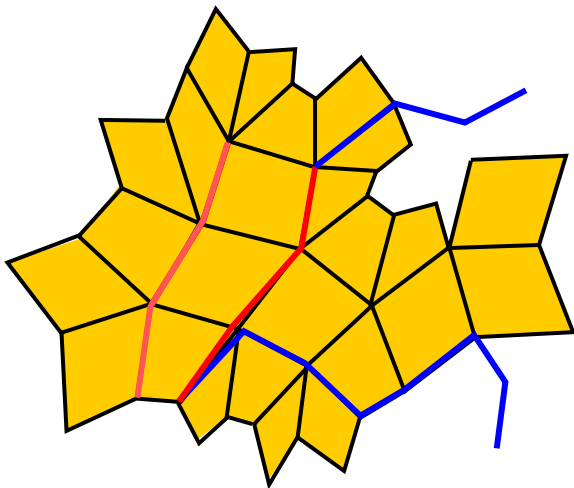
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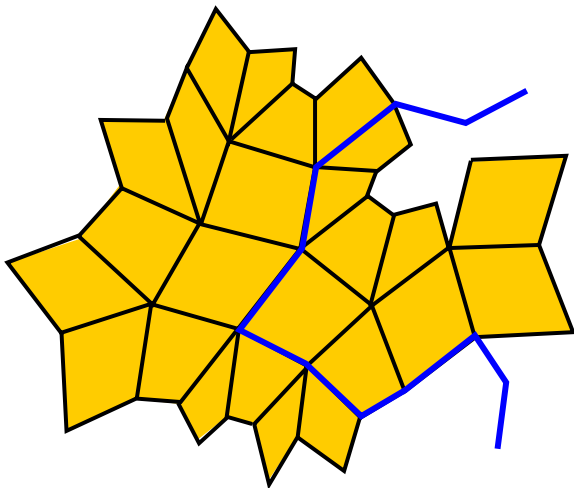
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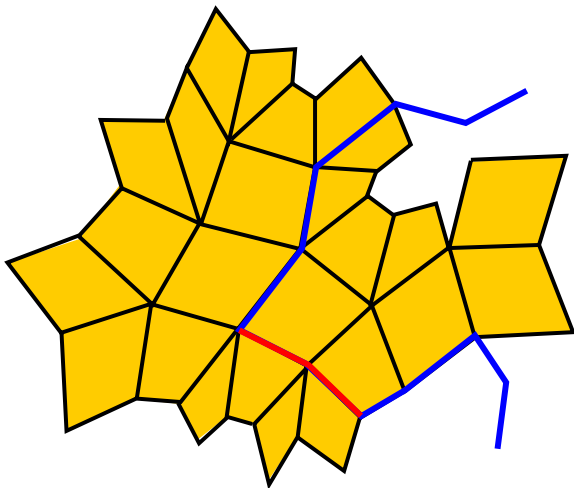
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Canonical representative

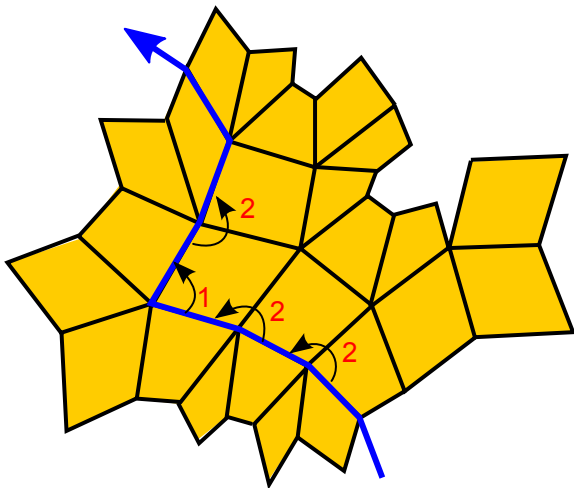
Hyp : All faces are quadrilaterals and all internal vertices have degree > 4 .



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Canonical representative

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We shorten by removing brackets and spurs and push to the right.

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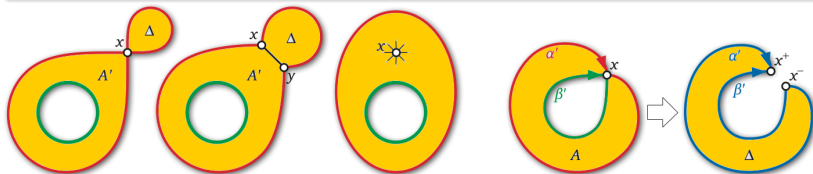
L. and Rivaud '12, Erickson and Whittlesey '13

After removing all spurs and brackets and pushing to the right as much as possible, we obtain a **canonical** representative. It can be computed in linear time.

Canonical representative

L. and Rivaud '12, Erickson and Whittlesey '13

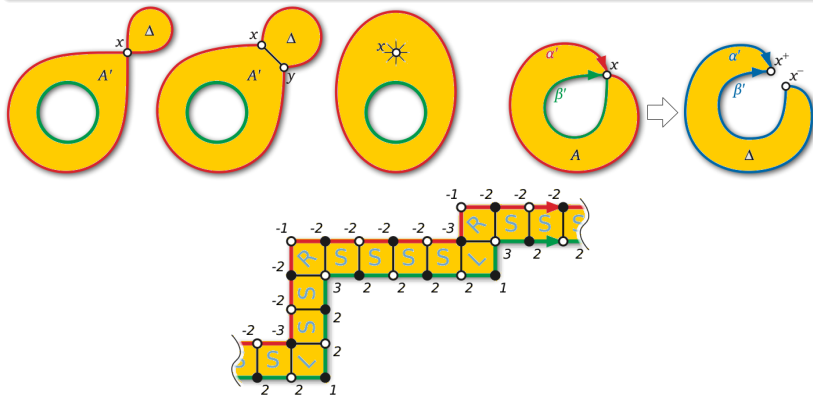
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One can decide if two curves are homotopic in linear time.

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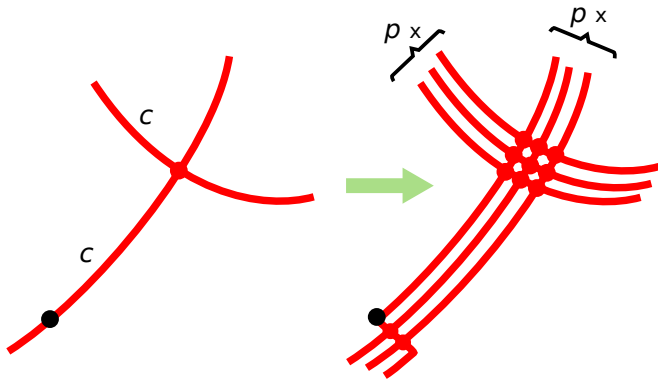
Corollary II

One can compute the primitive root of a curve in linear time.

Non-primitive curves

Lemma

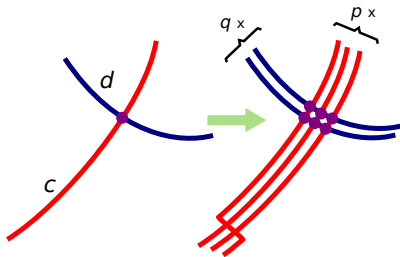
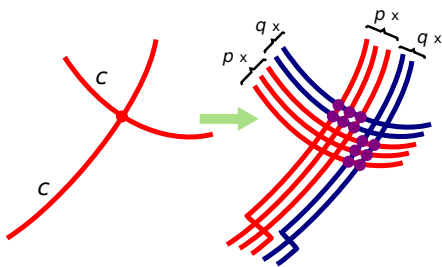
$$i(c^p) = p^2 \times i(c) + p - 1$$



Non-primitive curves

Lemma

$$\iota(c^p, d^q) = \begin{cases} 2pq \times \iota(c) & \text{if } c \sim d \text{ or } c \sim d^{-1}, \\ pq \times \iota(c, d) & \text{otherwise.} \end{cases}$$



The plan

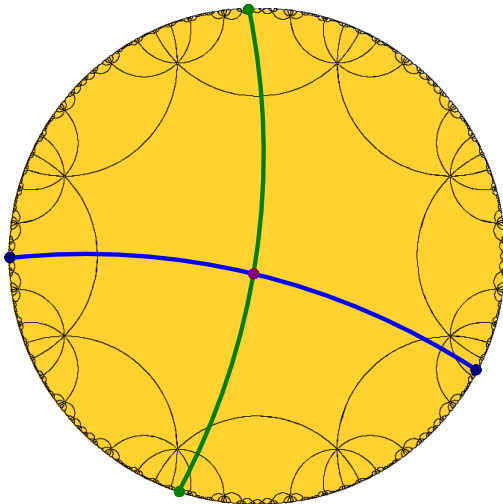
For a given curve c

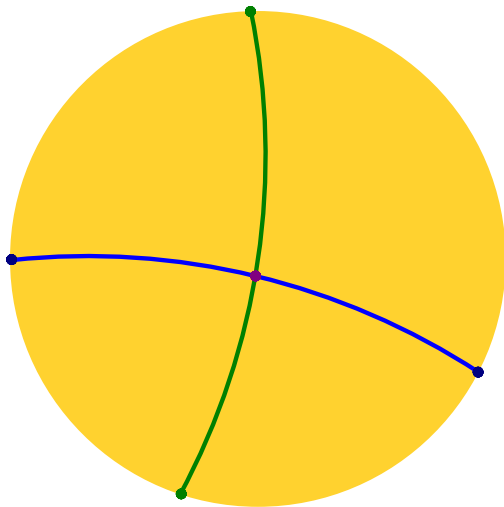
- 1 Determine the primitive root of c .
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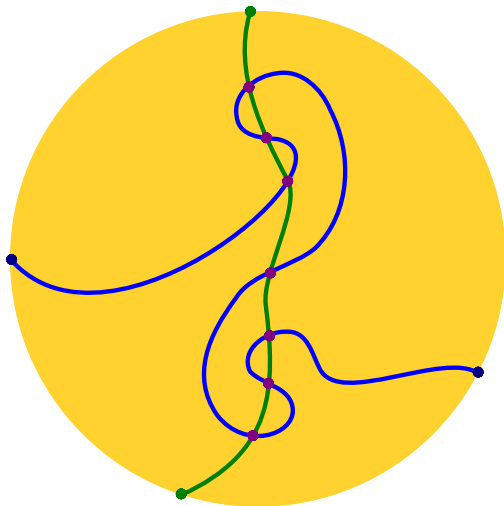
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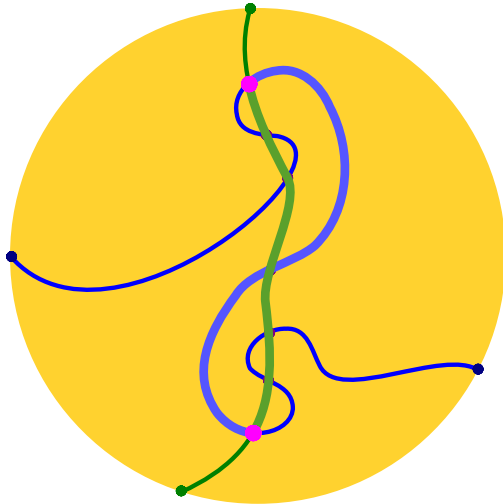
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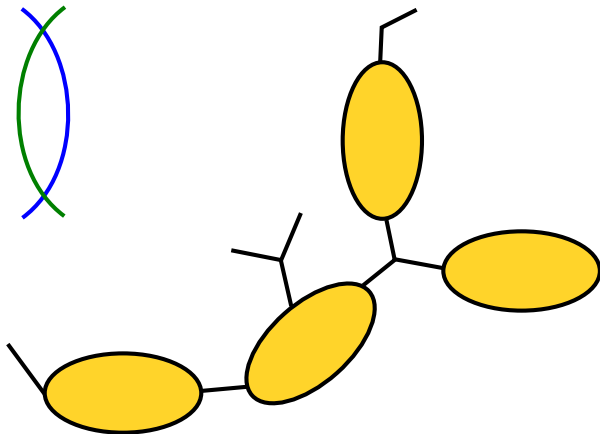


Double paths

A **double path** is a pair of homotopic paths.

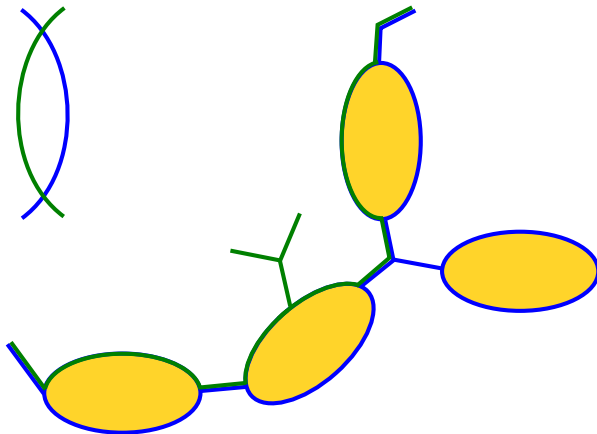
Lemma

Double paths in a canonical representative (resp. a geodesic) are **quasi-flat**.



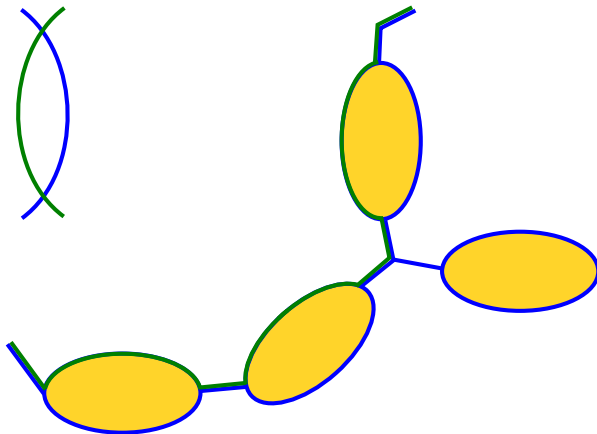
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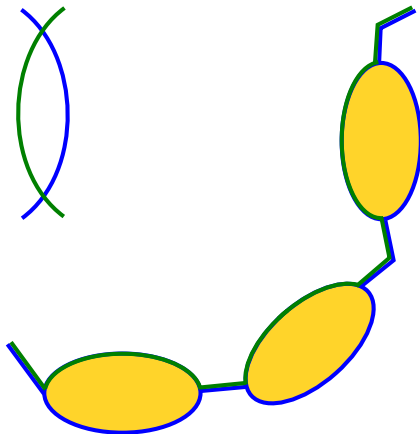
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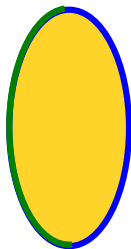
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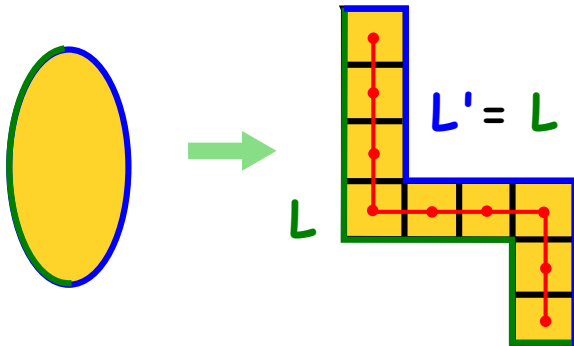
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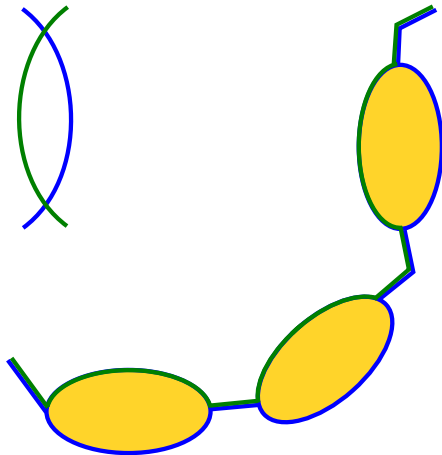
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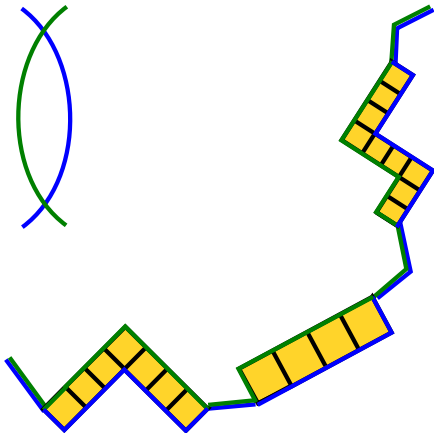
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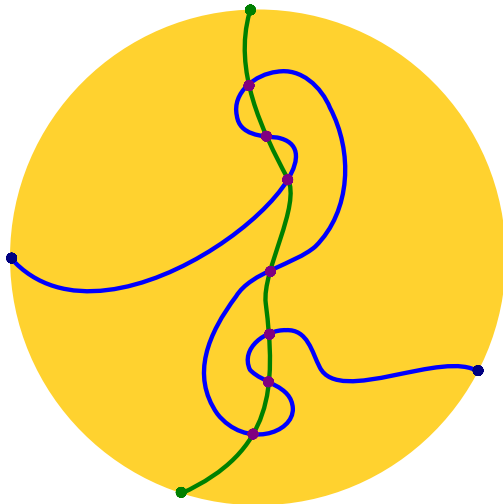
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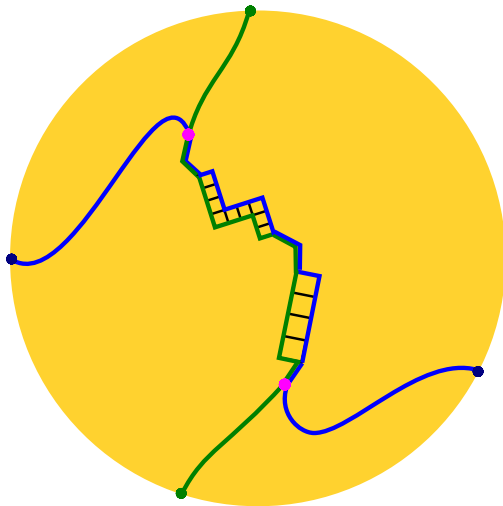
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The set of maximal double paths can be computed in quadratic time.

PROOF. A pair of indices (i, j) may occur at most twice in the set of maximal double paths. \square

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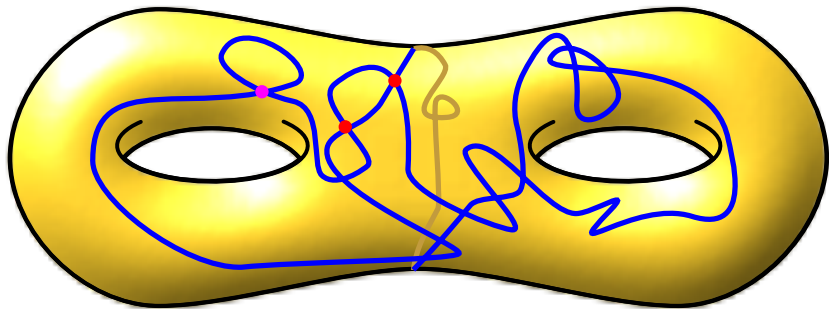
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Despré and L. '16

The geometric intersection number of one (two) curve(s) can be computed in quadratic time.

Computing an actual minimal configuration



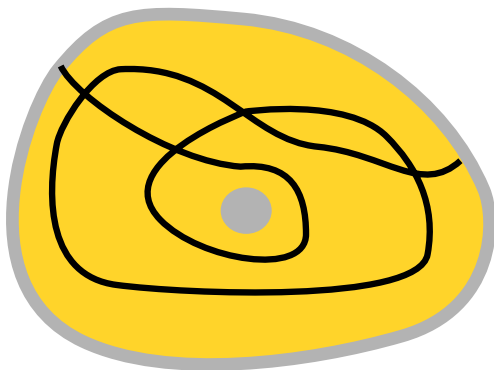
Hass and Scott '85

A curve with excess intersection has either a **monogon** or a **singular bigon**.

Computing an actual minimal configuration

Hass and Scott '85

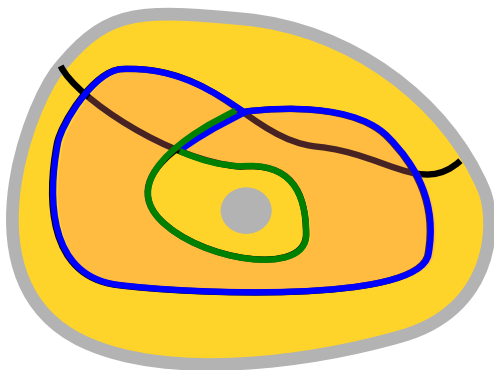
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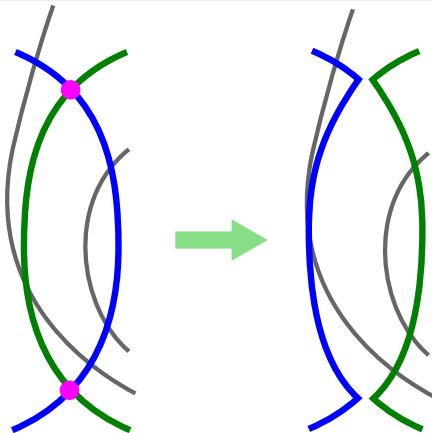
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Bigon swapping

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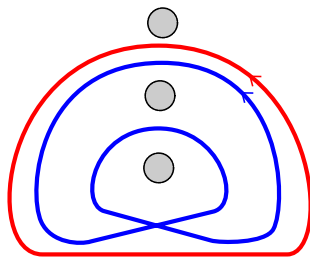
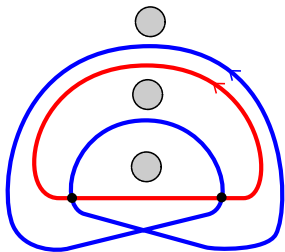
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Despré and L. '16

Given c , an homotopic immersion with a minimal number of intersections can be computed in $O(|c|^4)$ time.

Case of two curves



Open problems

- Propose an algorithm to compute a minimal immersion of two curves.

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- Propose an algorithm to compute a minimal immersion in hyperbolic configuration (cf. Hass and Scott '99).

Thank you for your
attention